



34

RACF

November 2003

In this issue

- 3 HLQ security check
 - 11 Business continuity and RACF
 - 20 PassTicket generator
 - 50 How RACF handles password
 - 62 RACF in focus – understanding OS/390 Unix security
 - 66 RACF news
-

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update

RACF Update

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RACF Update on-line

Code from *RACF Update*, and complete issues in Acrobat PDF format, can be downloaded from <http://www.xephon.com/racf>; you will need to supply a word from the printed issue.

Subscriptions and back-issues

A year's subscription to *RACF Update* (four quarterly issues) costs £190.00 in the UK; \$290.00 in the USA and Canada; £196.00 in Europe; £202.00 in Australasia and Japan; and £200.50 elsewhere. The price includes postage. Individual issues, starting with the August 1999 issue, are available separately to subscribers for £48.50 (\$72.75) each including postage.

Editor

Fiona Hewitt

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Printed in England.

HLQ security check

In a mainframe shop with z/OS, it's very important to check the 'dataset' type resource definitions – especially those belonging to the base software. Specifically, after the installation of a new version of the OS, you need to check that all target and distribution dataset prefixes are protected by RACF. It goes without saying that, because of the nature and use of the base software, all the prefixes examined refer only to datasets catalogued directly in the master catalog.

We wanted to avoid the situation whereby RACF, if it is active with the 'ProtectAll' option, might prevent a new address space from starting. To this end, I wrote the REXX code presented here.

The code should preferably be launched as batch, and produces a brief descriptive report of the installation. It can also be launched on the 'driving system' as long as it can access the SMP/E CSI that describes the 'target system'. The procedure will examine each dataset's high level qualifier to check that they are defined both in the Master Catalog and in a DDDEF of the target SMP/E.

In a correct installation, all the dataset prefixes that are defined in SMP/E should also be protected by RACF. The final list produced by this code will show which prefix is used by SMP/E and/or the Master Catalog and whether a RACF profile is missing or not.

QUAL1EXA REXX SAMPLE

```
/* REXX ----- */
/* REXX COMMAND : %QUAL1EXA MVS. GLOBAL. CSI_NAME TARGET_ZONE_NAME */
/* REXX ----- */
"PROF NOPREF"
PARSE UPPER ARG CSI _MVS TGT_ZON TRC_ACT
IF TRC_ACT = ' ON'
  THEN DO
```

```

TRACE ALL
MSG_STATUS = MSG("ON")
  END
  ELSE DO
  MSG_STATUS = MSG("OFF")
  TRACE OFF
  END

I F TGT_ZON = ' '
  THEN DO
  SAY ">>E>> TARGET ZONE NAME OMITTED "
  EXIT(15)
  END

AA = OUTTRAP(LINE.)
"LISTENT("CSI_MVS ") "
I F RC <> Ø
  THEN DO
  SAY ">>E>> 'MVS. CSI' NAME OMITTED "
  EXIT(25)
  END
LL = OUTTRAP(OFF)

/* - MASTER CATALOG ----- */
"FREE F(ICFQUAL1)"
"ALLOC F(ICFQUAL1) UNIT(VIO) RECFM(FB) LRECL(8)
  SPACE (5) TRACKS NEW DELETE"
I F RC > Ø THEN EXIT(35)
"EXECIO Ø DISK ICFQUAL1 (OPEN"

AA = OUTTRAP(LINE.)
"LISTNONVSAM "
I F RC > Ø THEN EXIT(45)
QUAL1_PREV = '_$_$_'

DO I =1 TO LINE. Ø

DSN_CAT = SUBSTR(LINE. I, 17, 1Ø)
POS_DOT = INDEX(DSN_CAT, '.' )

  I F POS_DOT = Ø
  THEN QUAL1_DSN = DSN_CAT
  ELSE DO
  QUAL1_LEN = POS_DOT - 1
  QUAL1_DSN = SUBSTR(DSN_CAT, Ø1, QUAL1_LEN)
  END

I F QUAL1_DSN <> QUAL1_PREV
  THEN DO
  QUAL1_PREV = QUAL1_DSN

```

```

        ICFQUAL1 = QUAL1_PREV
        PUSH ICFQUAL1
        "EXECIO 1 DISKW ICFQUAL1 "
        END

```

```

END

```

```

LL = OUTTRAP(OFF)

```

```

"EXECIO 0 DISKW ICFQUAL1 (FINIS)"

```

```

/* - DDDEF TARGET SMP/E DATASET ----- */

```

```

ML0000:

```

```

"FREE F(SMPLIST)"
"ALLOC F(SMPLIST) UNIT(VIO) RECFM(FBA) LRECL(121) BLKSIZE(7260)
    SPACE (5 5) TRACKS NEW DELETE"
"FREE F(SMPOUT)"
"ALLOC F(SMPOUT) DUMMY"
"FREE F(SMPLOG)"
"ALLOC F(SMPLOG) DUMMY"
"FREE F(SMPLOGA)"
"ALLOC F(SMPLOGA) DUMMY"
"FREE F(SMPRPT)"
"ALLOC F(SMPRPT) DUMMY"
"FREE F(SMPCSI)"
"ALLOC F(SMPCSI) SHRDSN("CSI_MVS")"
"FREE F(SMPCNTL)"
"ALLOC F(SMPCNTL) UNIT(VIO) RECFM(FB) LRECL(80) NEW DELETE"
RECNTL = "    SET BOUNDARY("TGT_ZON"). "
QUEUE RECNTL
RECNTL = "    LIST DDDEF.          "
QUEUE RECNTL
QUEUE ''
"EXECIO * DISKW SMPCNTL (FINIS)"
ADDRESS "LINKMVS" "GIMSMP"
    IF RC <> 0
        THEN DO
            SAY ">>E>> PGM : GIMSMP RETURNED ERROR "
            EXIT(55)
        END

```

```

/* ----- */

```

```

"FREE F(SMPWORK)"
"ALLOC F(SMPWORK) UNIT(VIO) RECFM(FB) LRECL(8)
    SPACE (5) TRACKS NEW DELETE"
    IF RC > 0 THEN EXIT(65)
"EXECIO 0 DISKW SMPWORK (OPEN)"
"EXECIO 0 DISKR SMPLIST (OPEN)"

```

```

DO FOREVER
"EXECIO 1 DISKR SMPLIST"
PULL REC_SMPLIST
IF RC <> 0 THEN LEAVE

IF SUBSTR(REC_SMPLIST, 12, 17) = ' DATASET      ='
    THEN DO
        SMPLIST_DSN = SUBSTR(REC_SMPLIST, 30, 10)
        POS_DOT     = INDEX(SMPLIST_DSN, '.')
        QUAL1_LEN   = POS_DOT - 1
        REC_SMPWORK = SUBSTR(SMPLIST_DSN, 01, QUAL1_LEN)
        PUSH REC_SMPWORK
        "EXECIO 1 DISKW SMPWORK "
        END
    END

END
"EXECIO 0 DISKW SMPWORK (FINIS"
"EXECIO 0 DISKR SMPLIST (FINIS"

/* - ICETOOL - START ----- */
"FREE F(SMPQUAL1)"
"ALLOC F(SMPQUAL1) UNIT(VIO) RECFM(FB) LRECL(8)
    SPACE (5) TRACKS NEW DELETE"
    IF RC > 0 THEN EXIT(75)
"FREE F(TOOLMSG)"
"ALLOC F(TOOLMSG) DUMMY"
"FREE F(DFSMSG)"
"ALLOC F(DFSMSG) DUMMY"
"FREE F(TOOLIN)"
"ALLOC F(TOOLIN) UNIT(VIO) RECFM(FB) LRECL(80) NEW DELETE"
REC_TOOLIN = " SELECT FROM(SMPWORK) TO(SMPQUAL1) ON(01, 8, CH) FIRST "
QUEUE REC_TOOLIN
QUEUE ''
"EXECIO * DISKW TOOLIN (FINIS"

ADDRESS "LINKMVS" "ICETOOL"
    IF RC <> 0 THEN EXIT(85)

/* - ICETOOL - END ----- */
"EXECIO 0 DISKR SMPQUAL1 (OPEN"
"EXECIO 0 DISKR ICFOQUAL1 (OPEN"
CALL READ_SMP
CALL READ_ICF
CALL PRINT_HEAD

DO FOREVER
END_FILE_ALL = SUBSTR(QUAL1_SMP, 1, 2)SUBSTR(QUAL1_ICF, 1, 2)

SELECT
WHEN END_FILE_ALL = '9999'

```

```

THEN LEAVE

WHEN QUAL1_SMP = QUAL1_ICF
  THEN DO
  CALL PRINT_PAGE
  CALL READ_SMP
  CALL READ_ICF
  END

WHEN QUAL1_ICF > QUAL1_SMP
  THEN DO
  CALL PRINT_PAGE
  CALL READ_SMP
  END

WHEN QUAL1_ICF < QUAL1_SMP
  THEN DO
  CALL PRINT_PAGE
  CALL READ_ICF
  END
END
END
SAY " * ----- * "

/* - END ----- */
EXIT(00)
/* ----- */

/* - INTERNAL ROUTINE - READ SMP ----- */
READ_SMP:
IF QUAL1_SMP <> '99SMP99'
  THEN DO
  "EXECIO 1 DISKR SMPQUAL1"
  IF RC > 0 THEN DO
  QUAL1_SMP = '99SMP99'
  "EXECIO 0 DISKR SMPQUAL1 (FINIS"
  END
  ELSE DO
  PULL QUAL1_SMP
  END
  END
RETURN

/* - INTERNAL ROUTINE - READ ICF ----- */
READ_ICF:
IF QUAL1_ICF <> '99ICF99'
  THEN DO
  "EXECIO 1 DISKR ICFQUAL1"

```

```

IF RC > 0 THEN DO
    QUAL1_ICF = '99ICF99'
    "EXECIO 0 DISKR ICFQUAL1 (FINIS"
        END
    ELSE DO
        PULL QUAL1_ICF
        END
    END
RETURN

/* - INTERNAL ROUTINE - PRINT PAGE ----- */
PRINT_PAGE:

PRINT_ICF = ''
PRINT_SMP = ''
PRINT_RACF = ''

SELECT

WHEN QUAL1_ICF = QUAL1_SMP
    THEN DO
        MY_QUAL1 = "('STRIP(QUAL1_ICF, 'T')'.*)"
        "LD DA"MY_QUAL1
        MY_RC = RC
        PRINT_ICF = QUAL1_ICF
        PRINT_SMP = QUAL1_SMP
        END

WHEN QUAL1_ICF < QUAL1_SMP
    THEN DO
        MY_QUAL1 = "('STRIP(QUAL1_SMP, 'T')'.*)"
        "LD DA"MY_QUAL1
        MY_RC = RC
        PRINT_ICF = QUAL1_ICF
        PRINT_SMP = '-----'
        END

WHEN QUAL1_ICF > QUAL1_SMP
    THEN DO
        MY_QUAL1 = "('STRIP(QUAL1_ICF, 'T')'.*)"
        "LD DA"MY_QUAL1
        MY_RC = RC
        PRINT_ICF = '-----'
        PRINT_SMP = QUAL1_SMP
        END

OTHERWISE NOP

END

```



```

IF MY_RC = 0 THEN PRINT_RACF = ' YES '
ELSE PRINT_RACF = ' NO '

```

```

P_A1 = ' * '
P_A2 = PRINT_ICF ' ' PRINT_SMP
P_A3 = ' ' PRINT_RACF
SAY P_A1 P_A2 P_A3 " " * "

```

```

RETURN

```

```

/* - INTERNAL ROUTINE - PRINT HEAD ----- */

```

```

PRINT_HEAD:

```

```

P_BL = ' '
W_DA = DATE(' S' )
X1 = SUBSTR(W_DA, 5, 2)
X2 = SUBSTR(W_DA, 7, 2)
X3 = SUBSTR(W_DA, 1, 4)
P_DA = X1"/"X2"/"X3
P_TI = TIME()
P_RA = SYSVAR("SYSLRACF")
P_IP = MVSVAR(' SYSNAME' )
P_MS = CSI_MVS
P_LL = P_MS P_BL P_BL
P_MI = SUBSTR(P_LL, 1, 44)

```

```

/* - CVT ----- */

```

```

CVT = C2D(STORAGE(10, 4))

```

```

/* - AMCBS ----- */

```

```

A_1 = C2D(STORAGE(D2X(CVT + 256), 4))

```

```

/* - ACB ----- */

```

```

A_2 = C2D(STORAGE(D2X(A_1 + 8), 4))

```

```

/* - CAXWA ----- */

```

```

A_3 = C2D(STORAGE(D2X(A_2 + 64), 4))

```

```

/* - MASTER CATALOG NAME ----- */

```

```

P_MT = STORAGE(D2X(A_3 + 52), 44)

```

```

SAY " * ----- * "

```

```

SAY " * " "TIME: " P_TI " DATE: " P_DA P_BL P_BL " * "

```

```

SAY " * " "SYSTEM ID: " P_IP P_BL P_BL P_BL P_BL " * "

```

```

SAY " * " "LVL/RACF : " P_RA P_BL P_BL P_BL P_BL " * "

```

```

SAY " * " "MVS MCAT : " P_MT " * "

```

```

SAY " * " "MVS CSI : " P_MI " * "

```

```

SAY " * ----- * "

```

```

SAY " *          MASTER          SMP/E          RACF          * "

```

```

SAY " * ----- * "

```

```

RETURN

```

```

-

```

SAMPLE JCL TO RUN QUAL1EXA

```
//..... JOB.....,.....,CLASS=.,MSGCLASS=.,REGION=0K
//** ----- **
//MYREXX EXEC PGM=IKJEFT01,DYNAMNBR=200,
//PARM='%QUAL1EXA.your.mvs_GLOBAL_CSI.your.mvs_target_zone'
//** ----- **
//SYSPRINT DD DUMMY
//SYSTSPRT DD SYSOUT=your.output_class
//SYSEXEC DD DISP=SHR,DSN=your.jcllib
//SYSTSIN DD DUMMY
```

PRINTOUT RESULT

```
* ----- *
* TIME:  xx:xx:xx  DATE:  xx/xx/xxxx  *
* SYSTEM ID:  IPOX  *
* LVL/RACF :  xxxx  *
* MVS MCAT :  CATALOG.MCAT.....  *
* MVS CSI :  .....GLOBAL.CSI  *
* ----- *
*          MASTER          SMP/E          RACF          *
* ----- *
*          AOP             AOP             YES           *
*          ASM             ASM             YES           *
*          ASU             ASU             YES           *
*          BFS             BFS             YES           *
*          CDS             CDS             NO            *
*          CEE             CEE             YES           *
*          -----        CICSSTS          NO            *
*          CSF             CSF             YES           *
*          DIT             DIT             YES           *
*          DRL             DRL             YES           *
*          DVG             DVG             YES           *
*          -----        IMS             YES           *
*          IMW             IMW             YES           *
*          IOA             IOA             YES           *
*          IOE             IOE             YES           *
*          ISF             ISF             YES           *
*          ISP             ISP             YES           *
*          -----        JAVA             YES           *
*          NETVIEW         NETVIEW         YES           *
*          OMVS            -----        YES           *
*          REXX            REXX             YES           *
*          -----        SMPZ13          YES           *
*          SOMMVS         SOMMVS          NO            *
*          SYS1            SYS1             YES           *
*          TCPIP          TCPIP             YES           *
*          TCPI VP        -----        YES           *
```

Business continuity and RACF

This article reviews the steps that you as RACF administrators can take, in conjunction with your technical support and business continuity departments, to ensure that your IT recovery site can continue processing with an adequate level of security. It discusses the options for RACF recovery from a variety of situations:

- Recovery at a third-party contingency site
- Recovery at your own separate contingency site
- LPAR-to-LPAR database recovery.

Let's face it: in a disaster situation, security is probably one of the last things people are really thinking about. Yet it can be one of the most important items required to make your recovery plans work properly. Not only that, it can also be the difference between the life and death of your company.

THE BASICS

Let's start with the simple stuff. What's your business continuity position in a disaster? Hiding under your desk in the foetal position, sucking your thumb, doesn't count – I'm talking about your company's position:

- Do you have a business continuity department? It could be called any of the following:
 - business continuity
 - business continuation

- business-as-usual
- contingency planning
- disaster recovery
- incident response.
- Is there an overall plan for recovery from a disaster situation?
 - are there specific, detailed technical recovery plans?
- Does the plan include an operating system recovery plan?
- Does the operating system recovery plan address RACF?
- Have the technical plans ever been tested?

If you have answered positively to all of the above, you're in business. If you haven't, you could be out of business if disaster strikes. It's that simple, folks.

For the purposes of this article, we'll assume that you have the items above. Now, how do you recover your security database? Well, that depends on your recovery options.

OPTION 1 – THIRD-PARTY RECOVERY SITE

The third-party recovery site option is very simple and straightforward. Back up the RACF database along with all of the other OS/390 system files. Recover them onto the DASD at the third-party site, and you're in business, right?

Well, no – wrong actually. You need to ask yourself some hard questions:

- How often is the RACF database backed up?
- What do you do if your back-ups run on Sunday at 04.00 and your system gets fried on Saturday at 22.00?
- Do you back up both the primary and secondary RACF databases?
- How long does it take to recover those files onto the third-party site's DASD?

- How do you handle DASD volume differences between sites?
- How do you handle IP address security between sites?

If at all possible, you should make sure that your RACF databases (primary and secondary) are backed up every night, preferably after the overnight production run. This back-up should then be sent along with your other off-site back-ups to a secure location. Please note, secure location does *not* mean in the same building as you made the back-ups, nor in a storage shed 20 feet away from the building where you made the back-ups. Try to ensure that your back-up tapes are stored at least one mile (1.6 km) away from your data centre. If you're in an area prone to earthquakes, make that a much longer distance, or transfer your files by secured high-speed data link.

Do you know how long it will take for your database to be recovered from tape onto the third-party system? Don't guess. When your company has a test of its recovery site, wangle a way to go with the rest of the team. You'll probably have to go anyway, just to fix the problems that invariably crop up. However, assuming that you've got all the back-up tapes, and all of them are readable (sometimes they're not, so a second set of back-ups is always a good idea), recovery of the operating system shouldn't take more than two to three hours.

Once the restore has been run and the operating system is up and running, you'll want to do two things. First, run any JCL you created before the test to update the DASDVOL settings within RACF, as well as any IP addresses you've hard-coded into the database. You should be able to get that information directly from your recovery site's technical support staff. Once this is done, run an IRRUT400 with INDEX and MAP to check the integrity of the database. The report is a bit long-winded, but it should give you some idea of any potential problems early on in the test.

As you progress through the recovery test, always make notes of any changes you make to the RACF database – access

permissions, ID sign-on problems, etc. You'll need this to enhance your recovery JCL, so that you can install all of those fixes early in the recovery process for the next test (or, of course, the real thing). Also, it would probably be a good thing to bring along a log of all of the updates you made to the RACF database in the days before the recovery test. It'll help you keep track of items that may have been missed if you're using a back-up from a few days past.

OPTION 2 – RECOVERY AT YOUR OWN SEPARATE CONTINGENCY SITE

Recovery at your own separate contingency site doesn't really differ much from the first option if your separate contingency site is a warm site (equipment there, but not loaded with operating system or software). If it's a hot site (fully loaded with programs, and just awaiting data), you need only restore the most recent RACF database from your production site and do some `RVARY SWITCH` and `RVARY ACTIVE` commands to replace the old database. You'll still need to stay on top of any operational security problems, of course, but if your company can afford its own contingency site, you'll save precious hours and minutes. That could easily make the difference between life and death for your organization – and that's no exaggeration.

OPTION 3 – LPAR-TO-LPAR DATABASE RECOVERY

One of my friends in technical support came up with this little jewel for LPAR-to-LPAR database recovery, so I can't claim full credit for the code presented here. However, I've found that it works very well, and it keeps the databases on separate LPARs in sync like a treat. It's also easy to run, and can be set up to execute the first part automatically on a daily basis (through automated production control software).

This option does require a couple of things, however. First, your systems must have a copy of Connect:Direct on board to use the JCL shown below. It'll probably also work with MQSeries or a simple ftp transfer as well, but these are options you'll need

to discuss with your technical support and network gurus. It also requires your LPARs to be in separate physical locations (ie city A and city B).

If you need to keep RACF databases in, let's say, a development and test LPAR in sync in a single box, this process will work. But the main focus of this article is recovery from a catastrophic loss of computer services. If you keep your production and recovery LPARs on the same piece of iron, you might as well tack a 'kick me' sign on your back. You're much safer having your recovery LPAR on a separate computer at a separate facility in a separate city. City-wide blackouts have been known to happen, so it's better to be safe than sorry. Besides, when the technical and operations staff do a test of their recovery plans, it gives you a good excuse to get out of town for a couple of days.

Anyway, back to the JCL. Run the first job on a daily basis in the early morning (my favourite time is 04:00) when your major production processing has finished. This job requires only one point of operator intervention, and that's just to ensure that Connect:Direct is up and running. After that point, the transfer goes on its merry way without the interference of mere mortals.

The second set of JCL is the one you submit yourself, from your own TSO library. Actually, before you submit the job, split your screen using F2 (if you're using IBM's PComm software) and have one frame showing the SDSF system log. You'll need it.

Submit the job, and then switch to the log. The mainframe will generate console messages asking for the RVAR Y ACTIVE and RVAR Y SWITCH commands. At the top of the SDSF log, type:

```
/nn, xxxxxxxx
```

where *nn* is the console message number and *xxxxxxx* is the password. You'll alternate between the VARY and SWITCH passwords, entering each one of them twice. The last step of the JCL performs a simple RVAR Y LIST, which you should use to ensure that your primary and back-up databases are in the right positions.

SOME FINAL COMMENTS ON BUSINESS CONTINUITY

- Some operations like to create 'recovery' userids for their contingency sites. You should avoid this like the plague! Those ids are usually of the 'super' variety, having access to everything. They also probably have OPERATIONS and SPECIAL capabilities as well. This may be a great shortcut, but it also short-circuits your overall security.
- Keep copies of your recovery plan off-site. Keep a copy at home. Keep it up to date. In the event of an emergency, you want to provide the fastest possible response to minimize recovery time. Digging around for hours in a storage facility, trying to find a particular plan, is a waste of time.
- Remember the old adage – 'Failure to plan on your part does not constitute a crisis on *my* part'. If other areas don't plan for disaster situations, that doesn't mean that you shouldn't as well. And if things do go wrong, you always have the option of saying those four little words that mean so much – "*I told you so!!!*"

JCL EXAMPLES

RACF database unload and transfer

```
//RACFDR JOB SYS1, ' SEND RACFDB TO DR' ,MSGCLASS=X,MSGLEVEL=(1,1),
//*      TYPRUN=HOLD,CLASS=A,NOTIFY=DFARMER,USER=PRODC TL
/*ROUTE XEQ PRD1
/*ROUTE PRINT PRD1
/*-----*
/* This job copies the RACF database to a DFDSS backup for *
/* transmission to a separate LPAR via Connect:Direct *
/* * *
/* This is the sequence of events: *
/* * *
/* Step1: Copy the RACF database to a separate file. *
/* * *
/* Step2: Send message to Console Operator to ensure Connect:Direct *
/*         is up and running. *
/* * *
/* Step3: Send the file to the receiving LPAR. *
/* * *
/*-----*
```



```

//CDWAIT EXEC PGM=CDWAIT, PARM=' 30: 30, YLLYYY'
//DMPUBLIB DD DSN=SYST. CD. PROCESS, DISP=SHR
//DMNETMAP DD DSN=SYST. CD. NETMAP, DISP=SHR
//DMMSGFIL DD DSN=SYST. CD. MSG, DISP=SHR
//NDMCMDS DD SYSOUT=*
//DMPRI NT DD SYSOUT=*
//SYSPRI NT DD SYSOUT=*
//SYSIN DD DSN=*&TEMP, DISP=(OLD, DELETE)
/*
//          ENDIF
//$J EXEC PGM=$A, COND=(0, LE), PARM=' RAC305270001000000000'

```

RACF restore job

```

//RACFREST JOB SYS1, ' RACF DB RESTORE', MSGCLASS=X, MSGLEVEL=(1, 1),
//      CLASS=A, NOTIFY=DFARMER, TYPRUN=HOLD
/*ROUTE XEQ BKP1
/*ROUTE PRI NT BKP1
/**-----*
/**   Thi s j ob restores the RACF database from a DFDSS backup      *
/**                                     *
/** Thi s i s the sequence of events:                                *
/**                                     *
/** Step1: I ssue RACF Command to l i st the state of the RACF datasets *
/**       before we start                                           *
/**                                     *
/** Step2: Wi l l swi tch to the backup RACF database, a repl y to   *
/**       confi rm the swi tch at the master consol e i s requi red. *
/**                                     *
/** Step3: Restores the Pri mary RACF database from a backup        *
/**                                     *
/** Step4: Acti vates the newl y restored RACF database, a repl y to *
/**       confi rm the acti vati on at the master consol e i s requi red. *
/**                                     *
/** Step5: Swi tches back to the Pri mary RACF database, a repl y to *
/**       confi rm the swi tch at the master consol e i s requi red. *
/**                                     *
/** Step6: Lock and copy the pri mary to the backup database       *
/**                                     *
/** Step7: Unl ocks the Pri mary after the copy has compl eted     *
/**                                     *
/** Step8: Acti vates the backup RACF database                     *
/**                                     *
/**-----*
/**       Query the status of RACF databases                        *
/**-----*
//STEP1 EXEC PGM=COMMAND
//INFILE DD *
RACF RVARY LI ST
/*

```

```

/*-----*
/*      Issue the swi tch command (from Primary database)      *
/*-----*
//      I F (STEP1. RC = 0) THEN
//STEP2 EXEC PGM=IKJEFT1A
//SYSTSPRT DD SYSOUT=*
//SYSTSI N DD *
RVARY SWI TCH DATASET(SYS1. RACF. RACFPRI M)
/*
//      ENDI F
/*-----*
/*      Restore the Primary RACF database from the backup      *
/*-----*
//      I F (STEP2. RC = 0) THEN
//STEP3 EXEC PGM=IRRUT400, PARM=' NOLOCKI NPUT, FREESPACE(20)'
//SYSPRI NT DD SYSOUT=*
//I NDD1 DD DSN=DSYS. PRD1BKUP. RACFPRI M, DI SP=OLD
//OUTDD1 DD DSN=SYS1. RACF. RACFPRI M, DI SP=OLD
//      ENDI F
/*-----*
/*      Acti vate the newly restored RACF database            *
/*-----*
//      I F (STEP3. RC = 0) THEN
//STEP4 EXEC PGM=IKJEFT1A
//SYSTSPRT DD SYSOUT=*
//SYSTSI N DD *
RVARY ACTI VE DATASET(SYS1. RACF. RACFPRI M)
/*
//      ENDI F
/*-----*
/*      Swi tch back to the Primary RACF database            *
/*-----*
//      I F (STEP4. RC = 0) THEN
//STEP5 EXEC PGM=IKJEFT1A
//SYSTSPRT DD SYSOUT=*
//SYSTSI N DD *
RVARY SWI TCH DATASET(SYS1. RACF. RACFBACK)
/*
//      ENDI F
/*-----*
/*      Lock and copy the primary to the backup database      *
/*-----*
//      I F (STEP5. RC = 0) THEN
//STEP6 EXEC PGM=IRRUT400, PARM=' LOCKI NPUT, FREESPACE(20)'
//SYSPRI NT DD SYSOUT=*
//I NDD1 DD DSN=SYS1. RACF. RACFPRI M, DI SP=OLD
//OUTDD1 DD DSN=SYS1. RACF. RACFBACK, DI SP=OLD
/*
//      ENDI F
/*-----*

```

```

/**      Unlock the primary      *
/**-----*
//      IF (STEP6.RC = 0) THEN
//STEP7 EXEC PGM=IRRUT400, PARM=' UNLOCKINPUT'
//SYSPRINT DD SYSOUT=*
//INDD1 DD DSN=SYS1.RACF.RACFPRI M, DISP=OLD
/*
//      ENDIF
/**-----*
/**      Activate the backup RACF database      *
/**-----*
//      IF (STEP7.RC = 0) THEN
//STEP8 EXEC PGM=IKJEFT1A
//SYSPRINT DD SYSOUT=*
//SYSDISP DD *
//RVARY ACTIVE DATASET(SYS1.RACF.RACFBACK)
/*
//      ENDIF
/**-----*
/**      Query the status of RACF databases      *
/**-----*
//      IF (STEP8.RC = 0) THEN
//STEP9 EXEC PGM=COMMAND
//INFILE DD *
//RVARY LIST
/*
//      ENDIF

```

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PassTicket generator

A RACF PassTicket is a one-time password which calculates a password based on the name of the application, the time, the userid, and a secret key. PassTickets are valid for approximately 10 minutes and generally can't be reused in that time. PassTickets are often used for Secure Single Sign-on procedures where a machine issues a log-on command to an application on behalf of an already authenticated user, without user intervention.

This article begins by explaining how PassTickets work and

examining some possible pitfalls. It ends by presenting the coding for a small ISPF PassTicket generator application.

PASSTICKETS PRIMER

Apart from the userid itself, a PassTicket has the following components.

The secret key is kept in the RACF database in a profile in the PTKTDATA class and consists of a 64-bit DES key that can be masked or encrypted if a cryptographic product like Integrated Cryptographic Service Facility (ICSF) is installed. If one of the systems is not RACF-controlled – an NT workstation, for instance – the place where the common secret, the DES key, is stored should be carefully checked. Once the secret key for an application is known, an intruder needs only a userid and the name of the application to generate a valid PassTicket.

The time is based on Greenwich Mean Time (GMT), nowadays called Universal Time Coordinates (UTC). For a PassTicket to function, the machines must be synchronized.

The application can be any VTAM application that passes the APPL= on the RACROUTE REQUEST=VERIFY,ENV=CREATE macro:

- TSO/E, for instance, does not, but there is a bypass hardcoded in the PassTicket algorithm.
- Pre-OS/390 V2R10 releases can use TSO<SMFID> as the application name.
- Starting from OS/390 V2R10, the VTAM Generic Resource Name (VGN) is used if defined in the GNAME= parameter in the TSOKEYxx member of the PARMLIB concatenation. Without a GNAME specification, TSO<SMFID> is used.
- For a batch job, the application name becomes MVS<SMF id>.

You should check the documentation of the application carefully before attempting to get PassTickets working. Sometimes a

bypass can be implemented in the RACROUTE REQUEST=VERIFY(X) pre-processing exit routine ICHRIX01.

More information on PassTickets can be found in *Security Server (RACF) Security Administrator's Guide* (SC28-1915). The complete calculation algorithm is described in *SecureWay Security Server RACF Macros and Interfaces* (SC28-1914). ICHRIX01 is described in *SecureWay Security Server RACF System Programmer's Guide* (SC28-1913).

USABILITY

In our shop, we use PassTickets for Secured Single Sign-on from the session manager to a number of applications both inside and outside the sysplex configuration that contains the network front-end machine. We also autoconnect trusted non-z/OS machines to the mainframe. Finally, we have an emergency system whereby after some operator commands an SOS RACF userid with extreme authorities gets resumed. The password can be generated with the application described below, which I also found very useful while implementing and testing PassTicket implementations. I've also seen PassTickets used to send a temporary password by e-mail to users who have forgotten their normal one – they then have about ten minutes to set a new password of their own.

RACF IMPLEMENTATION

In order to start using PassTickets for an application, the PTKTDATA class must be activated and RACLISTED, some profiles must be defined, and the synchronization of the clocks must be checked. See *SecureWay Security Server RACF Command Language Reference* (SC28-1919) for further details.

For the application TSOT and the user JEDSP00, a good start would be to execute the following RACF commands:

```
setropts cl assact(ptktdata)
setropts racl i st(ptktdata)
rdefi ne ptktdata tsot uacc(none) ssi gnon(keymasked(0123456789abcdef))
setropts refresh racl i st(ptktdata)
```

PTKTDATA PROFILES

PTKTDATA profiles can be defined in four different ways:

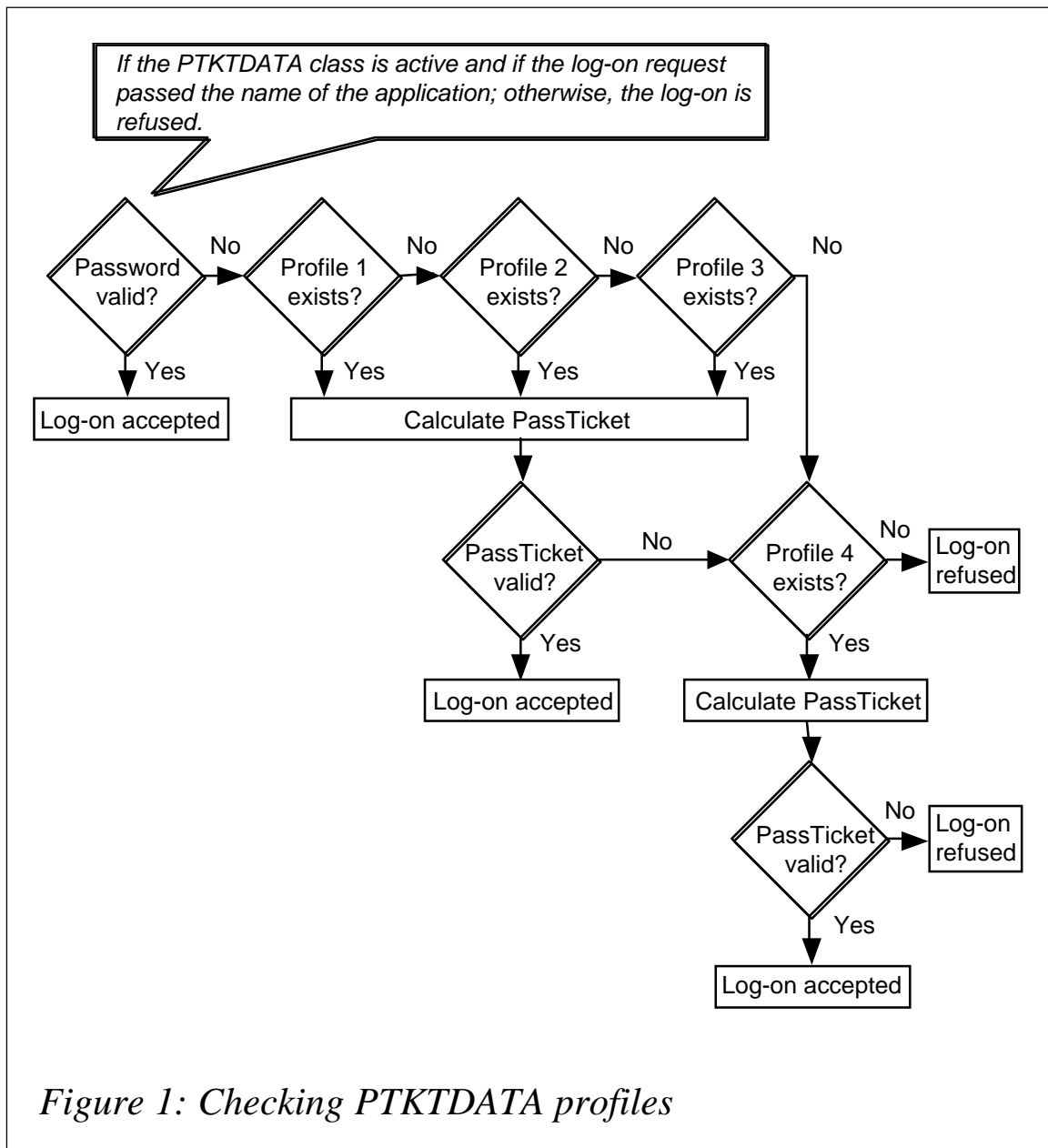
- 1 <application>.<group>.<user>
- 2 <application>.<user>
- 3 <application>.<group>
- 4 <application>

where <application> corresponds to the profile in the APPL class. The qualifier <group> is the actual connect group of the user. Some applications – such as TSO/E for instance – allow a user to pass a group in the log-on panel. The name is passed with the GROUP= parameter of the RACROUTE REQUEST=VERIFY,ENV=CREATE macro. This information is then stored in the ACEE field ACEEGRPN, and is the only one that the PassTicket algorithm will check – other groups to which the user is connected are ignored.

The first three profiles are said to be qualified, and only the last, non-qualified, profile is used to generate a PassTicket – we were very confused until we discovered this. For verification, on the other hand, if the string sent as a password doesn't correspond to the password in the RACF database, the PTKTDATA profiles are checked (see Figure 1).

Note that if profile 1 exists, profiles 2 and 3 aren't tested – this also confused us a lot.

We wanted to minimize the impact of a lost key in an insecure environment. We thought that giving different users different keys would ensure that only one could get compromised in the case of a breach. However, this didn't work out well in a RACF-RACF situation. Since the generation of the PassTicket uses only the <application> profile, this comes down to defining all users from RACF1 who want to access an application protected by RACF2 in one group in the RACF2 database. I don't think it's a good idea to create RACF groups solely for PassTicket verification in a RACF-RACF situation – the APPL class should take care of this.



THE GENERATOR

The PassTicket generator application consists of an Assembler program, a REXX program, and a small ISPF dialog. A REXX (POTEMKIN) program which displays the panel below calls the generator program:

```

-----
-----<Passticket Generator>-----
-----
  
```


VTAM Appl i cati on ===> TSOT
User I denti fi cati on ===> JEDSP00

PF3 to cancel
Enter to cal cul ate

The initial values for the two variables are taken from the user ISPF profile and are later restored there. After some formatting, the REXX program will call the POTEMKIN authorized Assembler program, which will do the following.

Before calculating a PassTicket, the program will check whether the user has SPECIAL RACF authority. If this is the case, a PassTicket is always generated. Otherwise, a comparison is made between the actual userid and the requested one, and the decision is made, based on RACF profiles in the PTKTDATA class (we thought it convenient to keep them there but the class can easily be changed if needed). The profiles are formatted in one of the two following ways:

- GENERATE.<application>.<group>
- GENERATE.<application>.<user>

READ is sufficient to generate a PassTicket for your own userid, but for another user UPDATE access is required. We check this for all the groups to which the target userid is connected.

Suppose that we have RACF 'systems' and 'security' groups, and we issue the following commands:

```
rdefi ne ptktdata generate. tsot. * uacc(none) owner(securi ty) +  
data(' Generate a PassTi cket for appl i cati on TSOT' )  
permi t generate. tsot. * cl ass(ptktdata) access(read) i d(*)  
permi t generate. tsot. * cl ass(ptktdata) access(update) i d(securi ty)  
rdefi ne ptktdata generate. tsot. systems uacc(none) owner(securi ty) +  
data(' Generate a PassTi cket for appl i cati on TSOT for a user connected to the group  
systems' )  
permi t generate. tsot. systems cl ass(ptktdata) access(update) i d(systems)  
setropts refresh racl i st(ptktdata)
```

With these rules in place, everybody can calculate a PassTicket for the application TSOT, but only for their own userid. This

could be an acceptable situation for a session manager to pass a log-on through to TSOT after the user has been authenticated. The 'security' group can generate a PassTicket for TSOT for everybody, and the 'systems' group can generate PassTickets for everybody who is connected to the 'systems' group. Note that there's no need for the KEY segment, since the profile isn't used by the PassTicket mechanism itself but only by the generator application.

If the user passes all the tests, the following message is displayed:

The PassTicket generated for user JEDSP00 and VTAM application TSOT is KDYMMU2G. This will be valid for 10 minutes and can only be used once.

Our operators use the application if they have to generate a PassTicket in an emergency. We therefore took care to generate as many clear messages as possible, so that no time would be lost in debugging. This made the coding rather lengthy.

THE DRIVING REXX PROGRAM (POTEMKIN)

```
/* REXX to display a panel with the appl and user id fields. It then
   tries to calculate a PassTicket.
*/

address "ISPEXEC"

"VGET (APPL USER) PROFILE"
user = strip(user)
appl = strip(appl)

"ADDP"
do forever
  "DISPLAY PANEL (PTKN000)"
  if strip(pf3) = 'END' ! appl = '' ! user = '' then leave
  if Get_Ptk() = 0 then leave
end

"REMP"
"VPUT (APPL USER) PROFILE"

exit
/* -----
Get_Ptk tries to calculate a PassTicket by calling the program
POTEMKIN. It analyses the return code (see ISPLIB) and
```

returns it to its caller.

```
----- */
Get_Ptk:

upper appl
upper user

l_user = length(user)
l_appl = length(appl)

user = left(user, 8)
appl = left(appl, 8)
parm = l_user user l_appl appl

address "TS0"

x = outtrap(' passticket. ')
"POTEMKIN" parm
retcode = rc
x = outtrap(' OFF' )

address "ISPEXEC"

if length(passticket. 1) > length(' PASSTICKET. 1') then do
  safrc = strip(left(passticket. 1, 7), 'L', '0')
  safrc = safrc !! substr(passticket. 1, 8, 1)
  x = substr(passticket. 1, 9, 8)
  racrc = strip(left(x, 7), 'L', '0')
  racrc = racrc !! substr(x, 8, 1)
  x = substr(passticket. 1, 17, 8)
  racrea = strip(left(x, 7), 'L', '0')
  racrea = racrea !! substr(passticket. 1, 8, 1)
  racrtype = strip(substr(passticket. 1, 25, 9))
  if racrtype = 'STAT' then ,
    msgpre = ' The PTKTDATA class or RACF is not active. '
  else msgpre = ''
end

select
  when retcode = 0 then do
    pkt = passticket. 1
    "SETMSG MSG(PTKN020)"
  end
  when retcode = 8 then "SETMSG MSG(PTKN000)"
  when retcode = 12 then "SETMSG MSG(PTKN001)"
  when retcode = 16 then "SETMSG MSG(PTKN002)"
  when retcode = 20 then "SETMSG MSG(PTKN003)"
  when retcode = 24 then "SETMSG MSG(PTKN004)"
  when retcode = 28 then "SETMSG MSG(PTKN005)"
  when retcode = 32 then "SETMSG MSG(PTKN006)"
```

```

when retcode = 36 then "SETMSG MSG(PTKN007)"
when retcode = 40 then "SETMSG MSG(PTKN008)"
when retcode = 44 then "SETMSG MSG(PTKN009)"
when retcode = 48 then do
  intyrc = strip(left(passticket.1, 7), 'L', '0')
  intyrc = intyrc !! substr(passticket.1, 8, 1)
  "SETMSG MSG(PTKN010)"
end
when retcode = 52 then "SETMSG MSG(PTKN011)"
when retcode = 56 then do
  cb = 'ASCB'
  "SETMSG MSG(PTKN012)"
end
when retcode = 60 then do
  cb = 'ASXB'
  "SETMSG MSG(PTKN012)"
end
when retcode = 64 then do
  cb = 'TCB'
  "SETMSG MSG(PTKN012)"
end
when retcode = 68 then "SETMSG MSG(PTKN013)"
when retcode = 72 then do
  cb = 'ACEE'
  "SETMSG MSG(PTKN012)"
end
when retcode = 76 then "SETMSG MSG(PTKN014)"
when retcode = 80 then "SETMSG MSG(PTKN003)"
when retcode = 84 then "SETMSG MSG(PTKN015)"
when retcode = 88 then "SETMSG MSG(PTKN016)"
when retcode = 92 then do
  cb = 'CVT'
  "SETMSG MSG(PTKN012)"
end
when retcode = 96 then do
  cb = 'RCVT'
  "SETMSG MSG(PTKN012)"
end
when retcode = 100 then "SETMSG MSG(PTKN017)"
when retcode = 104 then "SETMSG MSG(PTKN018)"
when passticket.1 = 'PASSTICKET.1' then "SETMSG MSG(PTKN019)"
otherwise nop
end

return(retcode)

```

THE ASSEMBLER COMMAND PROCESSOR (POTEMKIN)

In order to use the PassTicket callable service, a program must

run in key zero. This is a problem if you want to use an ISPF interface. What's more, REXX cannot trap output from a TPUT macro, ISPF services are difficult to use, the REXXATTACHMVS and LINKMVS environments abend if authorized functions are used, and the generated PassTicket is eight bytes, four bytes too long to be returned in register 15.

We decided to resolve these problems by writing a command processor. This allows the Assembler program to issue PUTLINE messages that can easily be trapped in REXX. An authorized command processor must be defined in the AUTHCMD macro of the IKJTSOxx member of the PARMLIB concatenation. A dynamic refresh is possible with the TSO/E PARMLIB UPDATE(xx) command. See also *TSO/E System Programming Command Reference* (SC28-1972) and *TSO/E Programming Services* (SC28-1971).

We link-edited the program in the LINKLIST concatenation with AC(1) as reentrant with AMODE 31 and RMODE ANY.

The RACROUTE REQUEST=EXTRACT allocates storage in subpool x'E5' which is private fetch protected storage. If you're using a FREEMAIN macro, no authorization is required. However, the STORAGE RELEASE macro requires storage key 0 or supervisor mode. See 'SecureWay Security Server RACF Macros and Interfaces' (SC28-1914) for information on the fields that can be retrieved.

The Assembler macro M#REGS equates the registers (RF is R15), and I imagine EYECATCH and AMODE31 are obvious.

```
TITLE ' *** POTEMKIN: CALCULATE A PASSTICKET          JANX
          DE DECKER ***'
*-----*
*  JAN DE DECKER JED: SP NV          JAN. DE. DECKER@TI SCALI . BE
*-----*
*
* NAME:      POTEMKIN TSO/E COMMAND PROCESSOR
*
* PARAMETERS: R1 -> CPPL
*             CPPLCBUF -> COMMAND BUFFER
*             COMMAND BUFFER: DISP PARAMETER LENGTH
*                               0  L(BUF)      2
```

```

*           2  N/A           2
*           4  COMMAND       8
*           D  L(USERID)    1
*           F  USERID       8
*          18  L(APPLID)     1
*          1A  APPLID        8
*

```

* PURPOSE: TSO COMMAND PROCESSOR TO CALCULATE A PASSWORD

* SYSTEM: OS/390 V2R10

* LINK: AMODE 31
RMODE ANY
REENTRANT
AC(1)
LINKLIST (REFRESH)

* INSTALL: LINK-EDIT INTO A LINKLIST LIBRARY AND REFRESH LLA
ADD TO IKJTSOXX AS AN AUTHORIZED COMMAND
RACF PTKTDATA CLASS ACTIVE AND GENERIC

* USE: SEE POTEMKIN REXX SAMPLE

* LOGIC: CALLED AS A TSO COMMAND WITH FIXED PARAMETERS
CHECKS - SEE RETURN CODES
IF CALLING USER HAS SPECIAL, CALCULATES PASSTICKET
IF CALLING USER IS TARGET USER, CHECKS READ ACCESS
OTHERWISE UPDATE. RULES THAT DETERMINE THE RIGHT
TO GENERATE A PASSTICKET ARE IN THE PTKTDATA CLASS
GENERATE. <APPL>. <USER OR A CONNECT GROUP>
IF SAF RETURNS ZERO, CALCULATES AND DISPLAYS
A PASSTICKET, OTHERWISE RETURNS

* RETURN CODES:	Ø ALL OK	PTKN020
*	8 NOT APF AUTHORIZED	PTKN000
*	12 STORAGE OBTAIN FAILURE	PTKN001
*	16 COMMAND BUFFER LENGTH ERROR	PTKN002
*	20 PTKTDATA CLASS OR RACF INACTIVE	PTKN003
*	24 LENGTH USERID > 8	PTKN004
*	28 LENGTH USERID = Ø	PTKN005
*	32 LENGTH APPL > 8	PTKN006
*	36 LENGTH APPL = Ø	PTKN007
*	40 USER IS UNDEFINED	PTKN008
*	44 USER IS REVOKED	PTKN009
*	48 ICHEINTY ERROR	PTKN010
*	52 ICHEINTY ERROR + PUTLINE ERROR	PTKN011
*	56 ASCB EYECATCHER ERROR	PTKN012
*	60 ASXB EYECATCHER ERROR	PTKN012
*	64 TCB EYECATCHER ERROR	PTKN012
*	68 NO ACEE FOUND	PTKN013

```

*          72 ACEE EYECATCHER ERROR          PTKN012
*          76 EMPTY ACEE USERID             PTKN014
*          80 RACROUTE ERROR                 PTKN003
*      84 NOT AUTHORIZED TO GENERATE PASSTICKET PTKN015
*          88 RACROUTE ERROR + PUTLINE ERROR PTKN016
*          92 CVT EYECATCHER ERROR          PTKN012
*          96 RCVT EYECATCHER ERROR         PTKN012
*      100 PASSTICKET GENERATION ERROR       PTKN017
*      104 PASSTICKET PUTLINE ERROR         PTKN018
*

```

```

* AUTHOR:      JAN          DATE: 06/2002
*

```

```

* SAMPLE:      N/A
*

```

```

* MODIFICATION:
*

```

```

*-----

```

```

          EJECT
POTEMKIN CSECT
POTEMKIN AMODE 31          31 BIT ADDRESSING
POTEMKIN RMODE ANY       PROGRAM CAN RESIDE ANYWHERE
          M#REGS .         REGISTER EQUATES
          PRINT NOGEN     DONT PRINT MACRO EXPANSIONS
          BAKR RE, 0      SAVE REGISTERS
          LR RC, RF       RC -> START OF POTEMKIN
          USING POTEMKIN, RC ADDRESS POTEMKIN WITH RC
          LR RA, R1       KEEP PARAMETER POINTER
          EYECATCH .     NAME, DATE & TIME OF ASSEMBLY
          AMODE31 .     CHANGES R0 AND R1
          EJECT
*

```

```

* START OF PROCESSING: CHECK AUTHORIZATION AND OBTAIN WORKING STORAGE
*

```

```

          TESTAUTH FCTN=1  ARE WE AUTHORIZED?
          LTR RF, RF       OK?
          BZ L0000        YES -> CONTINUE
          LA RF, 8        NO --> SET RC = 8
          PR .            AND RETURN TO CALLER
L0000 DS 0H
          STORAGE OBTAIN, ASK FOR STORAGE          X
          LENGTH=L_WORK, FOR THIS LENGTH        X
          COND=YES       CONDITIONALLY
          LTR RF, RF     STORAGE OBTAIN OK?
          BZ L0010      YES -> CONTINUE
          LA RF, 12     NO --> SET RC = 12
          PR .         AND RETURN TO CALLER
L0010 DS 0H
          LR RB, R1     RB -> WORK AREA
          LR R4, R1     R4 -> WORK AREA
          LR R6, R1     R6 -> WORK AREA

```

```

L    R7, =AL4(L_WORK)      R7 = L(WORK AREA)
XR   R5, R5                R5 = Ø
MVCL R6, R4                ZERO OUT WORK AREA
USING D_WORK, RB          RB ADDRESSES THE WORK AREA
EJECT

*
* INITIALIZE SOME FIELDS
*
LA   RD, SAVEAREA          RD -> SAVEAREA
MVC  SAVEAREA+4(4), =C' F1SA' LINKAGE STACK INDICATOR
MVC  RET_CODE, =F' Ø'      DEFAULT RETURN CODE = Ø
MVC  EXT_RA, =F' Ø'        RACROUTE EXTRACT RETURN AREA
EJECT

*
* ENTIRE PROGRAM: RB ADDRESSES OUR WORKAREA
*
RC   RC ADDRESSES OUR CSECT
RD   RD -> OUR SAVEAREA
* PREPARE THE MACRO AND I/O PARAMETER LIST ADDRESSES FOR PUTLINE
*
USING CPPL, RA            RA ADDRESSES THE CPPL
LA   R2, IOPLADS          R2 -> I/O PLIST ADDRESSES
USING IOPL, R2            R2 ADDRESSES THE IOPLADS
MVC  IOPLUPT, CPPLUPT     KEEP @(USER PROFILE TABLE)
MVC  IOPLECT, CPPLECT     KEEP @(ENV. CONTROL TABLE)
LA   R3, PUTL_ECB         R3 -> PUTLINE ECB
ST   R3, IOLECB           SET @(EVENT CONTROL BLOCK)
DROP R2                   FORGET THE IOPLADS
MVC  D_PUTL, S_PUTL       STATIC LIST MACRO TO DYNAMIC
EJECT

*
* CHECK LENGTH PARAMETER LIST IN COMMAND BUFFER
*
L    RA, CPPLCBUF         RA -> COMMAND BUFFER
DROP RA                   FORGET THE CPPL
USING P_LIST, RA          RA ADDRESSES THE COMMAND BUFFER
CLC  L_CMDBUF, =AL2(L_P_LIST) CHECK LENGTH PARAMETER LIST
BE   LØØ2Ø                OK -> CONTINUE
MVC  RET_CODE, =F' 16'    SET RETURN CODE
B    LØ31Ø                GO RETURN TO CALLER
EJECT

*
* CHECK IF THE CLASS PTKTDATA IS ACTIVE, IF NOT RETURN
*
LØØ2Ø DS ØH
MVC  D_RACSTA, S_RACSTA   COPY STATIC TO DYNAMIC MACRO
LA   R2, D_RACAUT         R2 -> RACROUTE MACRO
USING SAFR, R2            R2 ADDRESSES THE RACF PLIST
MVC  SAFPRRET, =F' Ø'     PRESET RACF RETURN CODE
MVC  SAFPRREA, =F' Ø'     PRESET RACF REASON CODE
DROP R2                   FORGET THE RACF PLIST

```


RACROUTE REQUEST=STAT,	REQUEST RACF INFORMATION	X
ENTRY=@_CDTENT,	WILL -> CDT ENTRY	X
RELEASE=2.6,	RACF RELEASE	X
WORKA=RACFWORK,	R7 -> RACF WORK AREA	X
MF=(E,D_RACSTA)	EXECUTE POINTING TO DYNAMIC	
LTR RF,RF	CLASS ACTIVE?	
BZ L0030	YES -> CONTINUE	
ST RF,X_SAF_RC	KEEP SAF RC	
LA R2,D_RACEXT	R2 -> RACROUTE PARAMETER LIST	
MVC RET_CODE,=F'20'	SET RETURN CODE	
MVC TEXT,BLANKS	BLANK OUT ERROR MESSAGE	
MVC RACRTYPE,=CL8'STAT'	SET RACROUTE TYPE	
B L0260	GO PRINT SAF AND RACF CODES	
EJECT		
*		
* THE LENGTH OF THE COMMAND BUFFER IS OK		
* RA ADDRESSES THE COMMAND BUFFER		
* COPY AND CHECK THE PARAMETER LIST		
*		
L0030 DS 0H		
MVC L_USERID,L_PARAM1	GET L(USERID)	
NI L_USERID,X'0F'	REMOVE FIRST HALF BYTE	
MVC USERID,PARAM1	GET USERID	
MVC L_APPLID,L_PARAM2	GET L(APPLID)	
NI L_APPLID,X'0F'	REMOVE FIRST HALF BYTE	
MVC APPLID,PARAM2	GET APPLID	
DROP RA	FORGET THE COMMAND BUFFER	
CLI L_USERID,X'08'	CHECK MAXIMUM LENGTH	
BNH L0040	NOT TOO HIGH, CONTINUE	
MVC RET_CODE,=F'24'	SET RETURN CODE	
B L0310	GO RETURN TO CALLER	
L0040 DS 0H	L(USERID) <= 8	
CLI L_USERID,X'00'	ZERO ?	
BNE L0050	NO --> CONTINUE	
MVC RET_CODE,=F'28'	SET RETURN CODE	
B L0310	GO RETURN TO CALLER	
L0050 DS 0H	0 < L(USERID) <=8	
NI L_APPLID,X'0F'	REMOVE FIRST HALF BYTE	
CLI L_APPLID,X'08'	CHECK MAXIMUM LENGTH	
BNH L0060	NO --> CONTINUE	
MVC RET_CODE,=F'32'	SET RETURN CODE	
B L0310	GO RETURN TO CALLER	
L0060 DS 0H	L(APPLID) <= 8	
CLI L_APPLID,X'00'	ZERO ?	
BNE L0070	NO --> CONTINUE	
MVC RET_CODE,=F'36'	SET RETURN CODE	
B L0310	GO RETURN TO CALLER	
L0070 DS 0H	0 < L(APPLID) <=8	
OC USERID,BLANKS	TO UPPER CASE	
OC APPLID,BLANKS		

```

EJECT
*
* TEST WHETHER OR NOT THE USER EXISTS AND IS NOT REVOKED
*
MVC L_WORK_R,=AL4(L'RACFWORK) LENGTH WORKAREA PREFIX
MVC D_INTY,S_INTY COPY STATIC TO DYNAMIC MACRO
ICHEINTY LOCATE, LOCATE A RACF PROFILE X
TYPE='USR', OF A USER ENTITY X
ENTRY=USER, POINTED TO BY R2 X
WKAREA=L_WORK_R, USE THIS WORK AREA X
TESTS=REVOKED, DO THE TEST AT LABEL REVOKED X
RELEASE=2.6, RACF RELEASE X
MF=(E,D_INTY) TARGET IS THE DYNAMIC ICHIEINTY
LTR RF,RF LOCATE OK?
BZ L0100 YES -> CONTINUE
C RF,=XL4'C' RETURN CODE 'C' = UNDEFINED
BNE L0080 NO --> DEFINED
MVC RET_CODE,=F'40' SET RETURN CODE
B L0310 GO RETURN TO CALLER
L0080 DS 0H USER IS DEFINED
C RF,=XL4'34' RETURN CODE '34' = REVOKED
BNE L0090 NOT UNDEFINED, NOT REVOKED
MVC RET_CODE,=F'44' SET RETURN CODE
B L0310 GO RETURN TO CALLER
EJECT
*
* WEIRD ICHIEINTY RETURN CODE: PRINT AND GO HOME
*
L0090 DS 0H
MVC RET_CODE,=F'48' SET RETURN CODE
ST RF,X_SAF_RC KEEP ICHIEINTY RC
LA RF,X_SAF_RC RF -> ICHIEINTY RC
ST RF,@_F_HEX STORE IN PLIST
LA RF,S_SAF_RC RF -> RECEIVE FIELD
ST RF,@_S_HEX STORE IN PLIST
LA R1,P_PRHGX R1 -> PARAMETER LIST
CALL F#PRHGX CALL FUNCTION TO PRINT FULLWORD
*
MVC PUTL_ECB,=F'0' ECB = 0
MVC TEXTADS,=H'12' L(OUTPUT LINE)
PUTLINE PARM=D_PUTL, PUTLINE LIST FORMAT X
OUTPUT=(TEXTADS,TERM,SINGLE,DATA), X
MF=(E,I OPLADS) PARAMETER LIST
LTR RF,RF PUTLINE OK?
BZ L0310 GO RETURN TO CALLER
MVC RET_CODE,=F'52' SET RETURN CODE
B L0310 GO RETURN TO CALLER
EJECT
*
* COMPARE THE PASSED USERID WITH THE USERID THAT ISSUED THE COMMAND

```

* IF THE USERIDS ARE THE SAME WE WILL CHECK READ ACCESS ON THE
 * GENERATE. <APPLICATION>. <USER OR CONNECT GROUP> RESOURCE
 * IN THE PKTDATA CLASS, OTHERWISE WE WILL CHECK UPDATE.
 * THE AUTHORITY REQUEST LEVEL WILL BE KEPT IN R6
 *

L0100	DS	0H	
	XR	R6, R6	R6 = 0
	LA	R6, 2	DEFAULT ATTR IS READ
	XR	R2, R2	R2 = 0
	USING	PSA, R2	R2 ADDRESSES THE PSA
	L	R2, PSAAOLD	R2 -> ASCB
	L	R3, PSATOLD	R3 -> TCB
	DROP	R2	FORGET THE PSA
	USING	ASCB, R2	R2 ADDRESSES THE ASCB
	CLC	=C' ASCB' , ASCBASCB	EYECATCHER OK?
	BE	L0110	YES -> CONTINUE
	MVC	RET_CODE, =F' 56'	SET RETURN CODE
	B	L0310	GO RETURN TO CALLER
L0110	DS	0H	R2 -> ASCB
	L	R2, ASCBASXB	R2 -> ASCB EXTENSION
	DROP	R2	FORGET THE ASCB
	USING	ASXB, R2	R2 ADDRESSES THE ASCB EXTENSION
	CLC	=C' ASXB' , ASXBASXB	EYECATCHER OK?
	BE	L0120	YES -> CONTINUE
	MVC	RET_CODE, =F' 60'	SET RETURN CODE
	B	L0310	GO RETURN TO CALLER
L0120	DS	0H	R2 -> ASXB
	L	R2, ASXBSENV	R2 -> ACEE
	DROP	R2	FORGET THE ASCB EXTENSION
	LTR	R2, R2	IS THERE A POINTER?
	BNZ	L0160	YES -> CONTINUE
L0130	DS	0H	NO --> CHECK TCB FOR POINTER
	LTR	R3, R3	SRB MODE? (PROBABLY YES)
	BZ	L0150	YES -> NO ACEE POINTER FOUND
	USING	TCB, R3	R3 ADDRESSES THE TCB
	CLC	=C' TCB' , TCBTCBID	EYECATCHER OK?
	BE	L0140	YES -> CONTINUE
	MVC	RET_CODE, =F' 64'	SET RETURN CODE
	B	L0310	GO RETURN TO CALLER
L0140	DS	0H	R2 -> TCB
	L	R2, TCBSENV	R2 -> ACEE
	DROP	R3	FORGET THE TCB
	LTR	R2, R2	IS THERE A POINTER
	BNZ	L0160	YES -> CONTINUE
L0150	DS	0H	NO ACEE COULD BE LOCATED
	MVC	RET_CODE, =F' 68'	SET RETURN CODE
	B	L0310	GO RETURN TO CALLER
L0160	DS	0H	R2 -> ACEE
	USING	ACEE, R2	R2 ADDRESSES THE ACEE
	CLC	=C' ACEE' , ACEEACEE	EYECATCHER OK?

BE L0170	YES -> CONTINUE	
MVC RET_CODE, =F' 72'	SET RETURN CODE	
B L0310	GO RETURN TO CALLER	
L0170 DS 0H	R2 -> ACEE	
CLC ACEEUSRI, =CL8' '	BLANK USERID?	
BNZ L0180	NO --> CONTINUE	
MVC RET_CODE, =F' 76'	SET RETURN CODE	
B L0310	GO RETURN TO CALLER	
L0180 DS 0H	NON-BLANK USERID IN THE ACEE	
CLC USERID, ACEEUSRI	ACEE USERID = PASSED USERID	
BE L0190	YES -> LEAVE READ ATTR	
LA R6, 4	NO --> SET ATTR TO UPDATE	
EJECT		
*		
* HAS THE COMMAND ISSUING USER THE RACF ATTRIBUTE SPECIAL?		
*		
L0190 DS 0H		
TM ACEEFLG1, ACEESPEC	CALLING USER HAS SPECIAL?	
* B0 L0270	YES -> CALCULATE PASSTICKET	
DROP R2	FORGET THE ACEE	
EJECT		
*		
* GET ALL THE GROUPS THE TARGET USER IS CONNECTED TO		
*		
MVC L_ENT_B, =H' 8'	LENGTH ENTITYX BUFFER	
MVC L_ENT_P, =H' 0'	LENGTH ENTITYX PROFILE NAME	
MVC ENT_PROF, BLANKS	BLANK OUT PROFILE	
XR R3, R3	R3 = 0	
IC R3, L_USERID	R3 = L(USERID)	
BCTR R3, 0	--R3 (FOR EX)	
LA R4, ENT_PROF	R5 -> PROFILE ENTITYX	
LA R5, USERID	R4 -> USERID	
EX R3, MVC1	MOVE IN APPLICATIONID	
*		
MVC D_RACEXT, S_RACEXT	COPY STATIC TO DYNAMIC MACRO	
LA R2, D_RACEXT	R2 -> RACROUTE MACRO	
USING SAFR, R2	R2 ADDRESSES THE RACF PLIST	
MVC SAFPRRET, =F' 0'	PRESET RACF RETURN CODE	
MVC SAFPRREA, =F' 0'	PRESET RACF REASON CODE	
DROP R2	FORGET THE RACF PLIST	
RACROUTE REQUEST=EXTRACT,	EXTRACT MACRO TYPE	X
TYPE=EXTRACT,	REQUIRE AN EXTRACT	X
ENTITYX=ENTITYX,	R5 -> ENTITY FIELD	X
FIELDS=EXFIELDS,	FIELDS TO EXTRACT	X
WORKA=RACFWORK,	R7 -> RACF WORK AREA	X
RELEASE=2.6,	RELEASE	X
MF=(E, D_RACEXT)	EXECUTE POINTING TO DYNAMIC	
LR R2, R1	R2 -> RESULT AREA	
LTR RF, RF	RACROUTE OK?	
BZ L0200	YES -> CONTINUE	

```

EJECT
*
* ERROR EXTRACTING THE CONNECT GROUPS OF THIS USER
*
      ST  RF, X_SAF_RC           KEEP SAF RC
      LA  R2, D_RACEXT          R2 -> RACROUTE PARAMETER LIST
      MVC RET_CODE, =F' 80'     SET RETURN CODE
      MVC TEXT, BLANKS          BLANK OUT ERROR MESSAGE
      MVC RACRTYPE, =CL8' EXTRACT' SET RACROUTE TYPE
      B   L0260                 GO PRINT SAF AND RACF CODES
      EJECT
*
* CONNECT GROUPS EXTRACT OK
*
L0200 DS 0H
      ST  R2, EXT_RA            KEEP ADDRESS RESULT AREA
      USING EXTWKEA, R2        R2 ADDRESSES FIXED RESULT AREA
      AH  R2, EXTWOFF          R2 -> VARIABLE RESULT AREA
      DROP R2                  FORGET THE FIXED RESULT AREA
*
* R2 -> +0 L(CONGRPCT DATA) = 4 L(CONNECT GROUP COUNT)
*      +4 CONGRPCT DATA       CONNECT GROUP COUNT
*      +8 L(CONGRPNM)          L(ALL CONNECT GROUPS)
*      +C L(CONGRPNM 1) = 8 L(FIRST CONNECT GROUP)
*      +10 CONGRPNM 1          FIRST CONNECT GROUP
*      +10 +L(CONGRPNM 1) L(CONGRPNM 2)
*
* R8 IS THE LOOP COUNTER : ALL THE CONNECTED GROUPS
* R9 -> CONNECTED GROUP ENTITY
*
      L   R8, 4(R2)             R8 = #(GROUPS)
      LA  R8, 1(R8)             + USERID
      LA  R9, 12(R2)           R9 -> FIRST GROUP DESCRIPTION
*
* PREPARE THE RACROUTE ON GENERATE. <APPL>. <USERID>
*      OR GENERATE. <APPL>. <GROUP> IN CLASS PTKTDATA
*
      MVC L_ENT_B, =H' 26'     MAXIMUM BUFFER LENGTH
      MVC L_ENT_P, =H' 0'     PROFILE LENGTH
*
L0210 DS 0H
      MVC ENT_PROF, BLANKS     BLANK OUT PROFILE
      MVC ENT_PRE, =CL9' GENERATE. ' MOVE IN PREFIX
      XR  R3, R3                R3 = 0
      IC  R3, L_APPLID          R3 = L(APPLID)
      BCTR R3, 0                --R3 (FOR EX)
      LA  R4, ENT_REM           R4 -> ENTITY PROFILE POSTFIX
      LA  R5, APPLID           R5 -> APPLICATION ID
      EX  R3, MVC1              MOVE IN APPLICATION ID
      AR  R4, R3                R4 -> LAST BYTE APPLICATION ID

```

```

LA R4, 1(R4)          R4 -> AFTER APPLICATION ID
MVI 0(R4), C' .'     INTER QUALIFIER
LA R4, 1(R4)          R4 -> AFTER ' .'
XR R3, R3             R3 = 0
IC R3, L_USERID      R3 = L(USERID)
BCTR R3, 0            --R3 (FOR EX)
LA R5, USERID        R5 -> USER ID
EX R3, MVC1           MOVE IN USERID
L0220 DS 0H
MVC D_RACAUT, S_RACAUT COPY STATI C TO DYNAMI C MACRO
LA R2, D_RACAUT      R2 -> RACROUTE MACRO
USING SAF, R2        R2 ADDRESSES THE RACF PLI ST
MVC SAFPRRET, =F' 0' PRESET RACF RETURN CODE
MVC SAFPRREA, =F' 0' PRESET RACF REASON CODE
DROP R2              FORGET THE RACF PLI ST
EJECT

*
* CALL RACF IN CLASS PTKTDATA WITH PROFILE GENERATE. <APPLID>. <USERID>
*
RACROUTE REQUEST=AUTH, REQUEST AUTHORITY X
ATTR=(R6), ACCESS REQUIRED X
ENTITX=(ENTITX, NONE), FOR THIS PROFILE X
RELEASE=2.6, RACF RELEASE X
WORKA=RACFWORK, WORKAREA X
MF=(E, D_RACAUT) EXECUTE FORMAT, TARGET IN R3
LTR RF, RF SAF OK?
BZ L0270 YES -> CONTINUE
ST RF, X_SAF_RC KEEP SAF RC
C RF, =F' 8' SAF CODE > 8?
BH L0250 NO --> GO PRINT SAF RACF CODES
EJECT

*
* LOOP FOR GROUPS
* WE ASSUME THAT THE LENGTH OF A GROUP NAME IS ALWAYS 8
*
LTR R8, R8 COUNTER 0?
BZ L0240 YES -> STOP
BCT R8, L0230 LOOP FOR ALL GROUPS
B L0240
L0230 DS 0H
MVC 0(8, R4), 4(R9) GET NEW GROUP
LA R9, 12(R9) R9 -> NEXT GROUP
B L0220
EJECT

*
* NO AUTHORIZATION TO CALCULATE A PASSTICKET
*
L0240 DS 0H
MVC RET_CODE, =F' 84' SET RETURN CODE
B L0310 GO RETURN TO CALLER

```

```

EJECT
*
* SAF RETURN CODE > 4 FROM RACROUTE REQUEST = AUTH
*
L0250 DS 0H
MVC RET_CODE, =F' 80' SET RETURN CODE
MVC TEXT, BLANKS BLANK OUT ERROR MESSAGE
MVC RACRTYPE, =CL8' AUTH' SET RACROUTE TYPE
LA R2, D_RACAUT R2 -> RACROUTE PARAMETER LI ST
EJECT
*
* WEIRD. RETURN ERROR SAF_RC RACF_RC RACF_REASON RACROUTE_TYPE
* R2 IS EXPECTED TO POINT TO THE RACROUTE PARAMETER LI ST
*
L0260 DS 0H
USING SAFP, R2 R2 ADDRESSES THE RACF PLI ST
LA RF, X_SAF_RC RF -> SAF RETURN CODE
ST RF, @_F_HEX STORE IN PLI ST
LA RF, S_SAF_RC RF -> RECEIVE FIELD
ST RF, @_S_HEX STORE IN PLI ST
LA R1, P_PRHGX R1 -> PARAMETER LI ST
CALL F#PRHGX CALL FUNCTION TO PRINT FULLWORD
*
LA RF, SAFPRRET RF -> RACF RETURN CODE
ST RF, @_F_HEX STORE IN PLI ST
LA RF, RACF_RC RF -> RECEIVE FIELD
ST RF, @_S_HEX STORE IN PLI ST
LA R1, P_PRHGX R1 -> PARAMETER LI ST
CALL F#PRHGX CALL FUNCTION TO PRINT FULLWORD
*
LA RF, SAFPRREA RF -> RACF REASON CODE
ST RF, @_F_HEX STORE IN PLI ST
LA RF, RACF_REA RF -> RECEIVE FIELD
ST RF, @_S_HEX STORE IN PLI ST
LA R1, P_PRHGX R1 -> PARAMETER LI ST
CALL F#PRHGX CALL FUNCTION TO PRINT FULLWORD
DROP R2 FORGET THE SAF PLI ST
*
MVC PUTL_ECB, =F' 0' ECB = 0
PUTLINE PARM=D_PUTL, PUTLINE LI ST FORMAT X
OUTPUT=(TEXTADS, TERM, SINGLE, DATA), X
MF=(E, IOPLADS) PARAMETER LI ST
LTR RF, RF PUTLINE OK?
BZ L0310 GO RETURN TO CALLER
MVC RET_CODE, =F' 88' SET RETURN CODE
B L0310 GO RETURN TO CALLER
EJECT
*
* CALCULATE THE PASSTICKET
*

```

L0270	DS	0H	
	XR	R2, R2	R2 = 0
	USING	PSA, R2	R2 ADDRESSES THE PSA
	L	R2, FLCCVT	R2 -> CVT
	DROP	R2	FORGET THE PSA
	USING	CVTMAP, R2	R2 ADDRESSES THE CVT
	CLC	=C' CVT' , CVTCVT	EYECATCHER OK?
	BE	L0280	YES -> CONTINUE
	MVC	RET_CODE, =F' 92'	SET RETURN CODE
	B	L0310	GO RETURN TO CALLER
L0280	DS	0H	R2 -> CVT
	L	R2, CVTRAC	R2 -> RCVT
	DROP	R2	FORGET THE CVT
	USING	RCVT, R2	R2 ADDRESSES THE RCVT
	CLC	=C' RCVT' , RCVTID	EYECATCHER OK?
	BE	L0290	YES -> CONTINUE
	MVC	RET_CODE, =F' 96'	SET RETURN CODE
	B	L0310	GO RETURN TO CALLER
L0290	DS	0H	R2 -> RCVT
	LA	RA, USER	RA -> USER FIELDS
	ST	RA, P_USER	STORE IN PLIST PTC
	LA	RA, APPL	RA -> APPL FIELDS
	ST	RA, P_APPL	STORE IN PLIST PTC
	MODESET	KEY=ZERO	STORAGE KEY ZERO
	LA	R1, P_PTC	R1 -> PLIST PASSTICKET CALC.
	L	RF, RCVTPTGN	RF -> PASSTICKET GENERATOR
	BALR	RE, RF	CALCULATE PASSTICKET
	DROP	R2	FORGET THE RCVT
	LR	RA, RF	KEEP RETURN CODE
	STM	R0, R1, TEXT	KEEP PASSTICKET
	MODESET	KEY=NZERO	TCB STORAGE KEY
	LTR	RA, RA	PASSTICKET SERVICE CALL OK?
	BZ	L0300	YES -> CONTINUE
	MVC	RET_CODE, =F' 100'	SET RETURN CODE
	B	L0310	GO RETURN TO CALLER
	EJECT		

*

* WRITE THE PASSTICKET TO THE TERMINAL USING PUTLINE

*

L0300	DS	0H	
	MVC	PUTL_ECB, =F' 0'	ECB = 0
	MVC	TEXTADS, =H' 12'	L(OUTPUT LINE)
	PUTLINE	PARM=D_PUTL,	PUTLINE LIST FORMAT
		OUTPUT=(TEXTADS, TERM, SINGLE, DATA),	X
		MF=(E, IOPLADS)	X
		PARAMETER LIST	
	LTR	RF, RF	PUTLINE OK?
	BZ	L0310	YES -> CONTINUE
	MVC	RET_CODE, =F' 104'	SET RETURN CODE
	B	L0310	GO RETURN TO CALLER
	EJECT		


```

*
* CLEAN UP THE ENVIRONMENT: STORAGE OBTAINED BY RACF EXTRACT
*
*
L0310 DS 0H
      L   RA, EXT_RA           RA -> RACROUTE EXTRACT RETURN
      LTR RA, RA              IS THERE ONE?
      BZ  L0320              NO --> CONTINUE, DON'T RELEASE
      USING EXTWKEA, RA      RA ADDRESSES THE RETURN AREA
      XR  R8, R8              R8 = 0
      IC  R8, EXTWSP         R8 = SUBPOOL ALLOCATED BY RACF
      XR  R9, R9              R9 = 0
      ICM R9, B'0111', EXTWLN R9 = L(EXTRACT RETURN AREA)
      DROP RA                FORGET THE RETURN AREA
      FREEMAIN R,           FREE UP THE RETURN AREA          X
          SP=(R8),          SUBPOOL NUMBER                X
          LV=(R9),          FOR THE GIVEN LENGTH          X
          A=(RA)            FROM THIS ADDRESS
L0320 DS 0H                WORK AREA CLEAN-UP
      L   RA, RET_CODE       RA = RETURN CODE
      STORAGE RELEASE,      FREE UP THE WORK AREA          X
          LENGTH=L_WORK,    FOR THE GIVEN LENGTH          X
          ADDR=(RB)         FROM THIS ADDRESS
      DROP RB                FORGET OUR WORK AREA
      EJECT

*
* END OF PROCESSING
*
THE_END DS 0H              MY ONLY FRIEND, THE END
      LR  RF, RA            LOAD RETURN CODE
      PR  .                 RETURN TO CALLER
      EJECT

*
* PR_HEX FUNCTION EXPECTS R5 TO POINT TO A FULL WORD AND R6 TO A CL8
* STRING. THE FULLWORD WILL BE PRINTED IN HEX FORMAT IN THE STRING
*
      BAKR RE, 0            SAVE REGISTERS
      PR  .                 RETURN TO CALLER
      EJECT

*
* EXECUTE TARGETS, VARIABLES AND CONSTANTS
*
MVC1 MVC 0(0, R4), 0(R5)    MOVE LAST QUALIFIER
BLANKS DC 133C' '
PTKTDATA DS 0F             CLASS ENTITY RACROUTE REQ=AUTH
      DC X'8'              SINGLE BYTE LENGTH
      DC CL8'PTKTDATA'     CLASS NAME
BIT0 DC X'80'             BIT IN ICHETEST MACRO
EXFIELDS DS 0F            FIELDS FOR RACROUTE REQ=EXTRACT
      DC A(2)              2 FIELDS

```

```

DC   CL8' CONGRPCT'      GROUP COUNT
DC   CL8' CONGRPNM'     GROUP NAMES
EJECT
*
* MACRO' S IN LIST FORMAT
*
S_PUTL  PUTLINE OUTPUT=(, TERM, SINGLE, DATA),          X
        MF=L
L_S_PUTL EQU *-S_PUTL      L(STATIC PUTLINE)
*
S_RACAUT RACROUTE REQUEST=AUTH,      REQUEST AUTHORITY          X
        ATTR=READ,          FOR READ ACCESS                      X
        CLASS=PTKTDATA,     IN THIS CLASS                        X
        ENTITYX=,          FOR THIS PROFILE                      X
        LOG=NOSTAT,        NO LOGGING                          X
        RELEASE=2.6,       RACF RELEASE                          X
        MF=L                LIST FORMAT
L_RACAUT EQU *-S_RACAUT    LENGTH OF RACROUTE MACRO
*
S_RACEXT RACROUTE REQUEST=EXTRACT,    EXTRACT MACRO TYPE        X
        TYPE=EXTRACT,      REQUIRE AN EXTRACT                  X
        CLASS=' USER' ,   USER PROFILE                        X
        ENTITYX=,          ENTITY FIELD                          X
        FIELDS,            FIELDS TO EXTRACT                    X
        WORKA=,            RACF WORK AREA                        X
        RELEASE=2.6,       RELEASE                              X
        MF=L                LIST_FORMAT
L_RACEXT EQU *-S_RACEXT    LENGTH OF RACROUTE EXTRACT MACRO
*
S_RACSTA RACROUTE REQUEST=STAT,      REQUEST RACF INFORMATION   X
        CLASS=' PTKTDATA' , CLASS NAME                          X
        ENTRY=,            RETURN ADDRESS -> CDT ENTRY          X
        RELEASE=2.6,       RACF RELEASE                          X
        WORKA=,            RACF WORK AREA                        X
        MF=L                LIST_FORMAT
L_RACSTA EQU *-S_RACSTA    LENGTH OF RACROUTE STAT MACRO
*
S_INTY  ICHEINTY LOCATE,             LOCATE A PROFILE          X
        TYPE='USR' ,       OF A USER ENTITY                    X
        ENTRY=,            POINTED TO BY R2                      X
        WKAREA=,          USE THIS WORK AREA                    X
        RELEASE=2.6,       RACF RELEASE                          X
        TESTS=,           DO THE TEST AT LABEL REVOKED         X
        MF=L
L_INTY  EQU *-S_INTY          LENGTH OF ICHEINTY MACRO
*
REVOKED ICHETEST FIELD=FLAG4,        X
        FLDATA=(1, BI TØ),          X
        COND=ZEROS
EJECT

```

```

*
* LITERAL POOL
*
      LTORG
      EJECT
*
* EQUATES
*
      PRINT NOGEN
*
* DYNAMIC WORK AREA DSECT
*
D_WORK DSECT
SAVEAREA DS 18F          SAVE AREA
RET_CODE DS F           RETURN CODE
X_SAF_RC DS F           SAF RC
EXT_RA DS F             RACROUTE EXTRACT RETURN AREA
@_UPT DS F              KEEPS @(UPT)
@_ECT DS F              KEEPS @(ECT)
@_CDTENT DS F          WILL KEEP @(CDT ENTRY PTKTDATA)
*
* PARAMETERS FOR THE PUTLINE MACRO
*
TEXTADS DS H            TEXTBUFFER LENGTH
      DS H              RESERVED
TEXT DS 0CL32          PASSTICKET OR RC'S
S_SAF_RC DS CL8        SAF RC STRING
RACF_RC DS CL8         RACF RC STRING
RACF_REA DS CL8        RACF REASON CODE STRING
RACRTYPE DS CL8        RACROUTE REQUEST=
*
IOPLADS DS 4F          PUTLINE PARAMETER LIST
PUTL_ECB DS F          PUTLINE ECB
D_PUTL DS XL(L_S_PUTL) PUTLINE DYNAMIC FORMAT
      EJECT
*
* DYNAMIC RACF FIELDS
*
ENTITX DS 0F           RACROUTE ENTITX
L_ENT_B DS H           BUFFER LENGTH
L_ENT_P DS H           PROFILE LENGTH IF KNOWN
ENT_PROF DS 0CL26      PROFILE
ENT_PRE DS CL9         'GENERATE.'
ENT_REM DS CL17        REMAINDER OF PROFILE
*
D_RACAUT DS XL(L_RACAUT) DYNAMIC RACROUTE REQUEST=AUTH
D_RACEXT DS XL(L_RACEXT) DYNAMIC RACROUTE REQUEST=EXTRACT
D_RACSTA DS XL(L_RACSTA) DYNAMIC RACROUTE REQUEST=STAT
D_INTY DS XL(L_INTY)   DYNAMIC ICHEINTY
*

```

* PARAMETERS FOR THE PASSTICKET CALCULATOR SERVICE ROUTINE

*

P_PTC	DS	ØF	PLIST PASSTICKET CALCULATOR
P_USER	DS	F	-> USER FIELDS
P_APPL	DS	F	-> APPLICATION FIELDS

*

USER	DS	ØH	KEEP INFORMATION TOGETHER
L_USERID	DS	X	L(USERID)
USERID	DS	CL8	USERID PADDED WITH BLANKS
APPL	DS	ØH	KEEP INFORMATION TOGETHER
L_APPLID	DS	X	L(APPLID)
APPLID	DS	CL8	APPLID PADDED WITH BLANKS

*

* PARAMETERS FOR THE PRINT A HEXADECIMAL NUMBER FUNCTION

*

P_PRHEX	DS	ØF	START OF PARAMETER LIST
@_F_HEX	DS	F	-> FULLWORD
@_S_HEX	DS	F	-> PRINTABLE STRING

*

* LAST ADDRESSABLE LABEL

*

L_WORK_R	DS	F	L(RACF WORK AREA) - ICHENTY
RACFWORK	DS	XL4Ø96	RACF WORK AREA

*

* END OF WORK AREA

*

L_WORK	EQU	*-D_WORK	LENGTH OF THE WORKAREA
		EJECT	

*

* PARAMETER LIST: THIS IS THE TSO/E COMMAND BUFFER

*

SINCE THE PARAMETER LIST IS CONSTRUCTED BY THE

*

CALLING REXX, WE DON'T TEST TOO MUCH IN THE PROGRAM

*

P_LIST	DSECT	.	PARAMETER LIST
L_CMDBUF	DS	H	L(COMMAND BUFFER)
		DS H	SEE TSO/E PROGRAMMING SERVICES
COMMAND	DS	CL8	OUR TMP
		DS X	BLANK
L_PARAM1	DS	X	L(USERID)
		DS X	BLANK
PARAM1	DS	CL8	USERID
		DS X	BLANK
L_PARAM2	DS	X	L(APPL)
		DS X	BLANK
PARAM2	DS	CL8	APPLID
L_P_LIST	EQU	*-P_LIST	L(PARAMETER LIST)
		EJECT	

*

* SYSTEM DSECTS

*

```

PRINT GEN
CVT DSECT=YES, PREFIX=YES COMMUNICATION VECTOR TABLE
EJECT
IHAPSA DSECT=YES PREFIXED STORAGE AREA
EJECT
ICHSAFP . RACROUTE PARAMETER LIST
EJECT
ICHPRCVT . RACF CVT
EJECT
IKJCPPL . COMMAND PROCESSOR PLIST
EJECT
IKJIOPL . I/O PARAMETER LIST ADDRESSES
EJECT
IHAASCB . ADDRESS SPACE CONTROL BLOCK
EJECT
IHAASXB . ASCB EXTENSION
EJECT
IHAACEE . ACCESS CONTROL ENVIRONMENT EL.
EJECT
IKJTCB . TASK CONTROL BLOCK
EJECT
IRRPRTW . RACROUTE EXTRACT RESULT AREA
END
EJECT

```

```

*-----
TITLE '*** JED: SP REPORTS: PRINT A HEXADECIMAL NUMBER JANX
DE DECKER ***'

```

```

*-----
* JED: SP JAN. DE. DECKER@TISCALI.BE
*-----

```

```

* NAME: F#PRHEX
*
* PURPOSE: PRINT A HEXADECIMAL NUMBER
*
* PARAMETERS: R1 -> @(FULLWORD)
                @(PRINT FULLWORD)
*
* LINK: CAN BE LINKED REENTRANT
*
* SYSTEM: OS/390 V2R10
*
* MODIFICATION:
*
*-----

```

```

F#PRHEX CSECT
F#PRHEX AMODE 31 31 BIT ADDRESSING
F#PRHEX RMODE ANY PROGRAM CAN RESIDE ANYWHERE
M#REGS . REGISTER EQUATES

```

```

EJECT
PRINT NOGEN
BAKR RE, Ø
LR RC, RF
USING F#PRHEX, RC
LR RA, R1
EYECATCH .
AMODE31 .
EJECT

*
* START PROCESSING
*
      LM    R5, R6, Ø(RA)          R5 -> FULLWORD
R6 -> PRINTABLE FULLWORD
      XR    R2, R2                R2 = Ø
      LA    R4, 4                  R4 = 4
LØØØØ DS   ØH
      IC    R2, Ø(R5)             LOAD BYTE INTO R2
      SRL   R2, 4                  SHIFT 4 BITS TO THE RIGHT
      LA    R3, 2                  R3 = 1
LØØ1Ø DS   ØH
      CH    R2, =H' 1Ø'           R2 >= X' A' ?
      BL    LØØ2Ø                 NO --> GO ADD FØ
      AH    R2, =H' 183'         YES -> ADD CØ
      B     LØØ3Ø                 NEXT HALFBYTE
LØØ2Ø DS   ØH
      AH    R2, =H' 24Ø'         ADD X' FØ'
LØØ3Ø DS   ØH
      STC   R2, Ø(R6)            STORE IN RECEIVE FIELD
      LA    R6, 1(R6)           POINT TO NEXT BYTE IN RECEIVE F
      IC    R2, Ø(R5)           TAKE THE SAME BYTE
      N     R2, =X' ØØØØØØØF'   MAKE DI SAPPEAR THE FIRST HALFB.
      BCT   R3, LØØ1Ø           AND JUMP 1 TIME TO LØØ1Ø
      LA    R5, 1(R5)           POINT TO NEXT FULLWORD BYTE
      BCT   R4, LØØØØ           JUMP 3 TIMES
      EJECT

*
* END OF PROCESSING
*
THE_END DS   ØH
      XR    RF, RF
      PR    .
      EJECT

*
* LITERAL POOL
*
      LTORG

*
      END

```

THE ISPF PANEL (PTKN000)

```
)ATTR DEFAULT(1#{})
E TYPE(PT)
] TYPE(NT)
{ TYPE(NEF) PADC(USER)
} TYPE(RP)
)BODY EXPAND($$) WINDOW(43, 11)
E-$-$-
E-$-$-<Passticket Generator>-$-$-
E-$-$-
E
] VTAM Application ==> {APPL ]
] User Identification ==> {USER ]
]
]
} PF3 to cancel
} Enter to calculate
]
)INIT
. CURSOR = APPL
)PROC
&PF3 = . RESP
)END
```

THE ISPF MESSAGES

PTKN00 member

```
PTKN000 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEMKIN Return code &RETCODE: ' +
' The POTEMKIN program is not authorized. Check APF authorization, ' +
' and the settings of IKJTS0xx. POTEMKIN must be in the AUTHCMD' +
' list to be invoked from ISPF. '
```

```
PTKN001 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEMKIN Return code &RETCODE: ' +
' The STORAGE OBTAIN SVC failed. '
```

```
PTKN002 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEMKIN Return code &RETCODE: ' +
' The command buffer passed to the generator has the wrong length. '
```

```
PTKN003 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEMKIN Return code &RETCODE: ' +
' &MSGPRE The RACROUTE' +
' REQUEST=STAT macro returned &SAFRC with a RACF return code' +
' &RACFRC and a RACF reason code &RACFREA. . '
```

PTKN004 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEMKIN Return code &RETCODE: ' +
' The passed user identification is longer than 8 characters. '

PTKN005 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEMKIN Return code &RETCODE: ' +
' The passed user identification has a zero length. '

PTKN006 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEMKIN Return code &RETCODE: ' +
' The passed VTAM application name is longer than 8 characters. '

PTKN007 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEMKIN Return code &RETCODE: ' +
' The passed VTAM application name has a zero length. '

PTKN008 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEMKIN Return code &RETCODE: ' +
' The user id &USER is unknown to RACF. '

PTKN009 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEMKIN Return code &RETCODE: ' +
' The user id &USER is REVOKED. '

PTKN01 member

PTKN010 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEMKIN Return code &RETCODE: ' +
' The ICHENTY macro returned unexpectedly (hex): &INTYRC. . '

PTKN011 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEMKIN Return code &RETCODE: ' +
' The ICHENTY macro returned an unexpected return code but' +
' unfortunately also the PUTLINE macro used to return this to' +
' the driving REXX program failed. '

PTKN012 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEMKIN Return code &RETCODE: ' +
' There was an error in the eyecatcher of the control block: &CB. . '

PTKN013 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEMKIN Return code &RETCODE: ' +
' An ACEE could not be located. '

PTKN014 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEMKIN Return code &RETCODE: ' +
' The ACEE user id contained only blanks. '

PTKN015 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEMKIN Return code &RETCODE: ' +

' You are not authorized by the RACF profiles GENERATE. &APPL ' +
' postfixed by &USER or a CONNECT GROUP in the class PTKTDATA ' +
' to generate a PassTicket for the user &USER and ' +
' the VTAM application &APPL. . ' +

PTKN016 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEKIN Return code &RETCODE: ' +
' There was a non-zero return code from a RACROUTE macro. ' +
' Unfortunately there was also an error while trying to ' +
' print the SAF and RACF return and reason codes using ' +
' the TSO/E PUTLINE macro. ' +

PTKN017 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEKIN Return code 56: ' +
' The PassTicket generation failed for user &USER and VTAM ' +
' application &APPL. . ' +

PTKN018 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' POTEKIN Return code &RETCODE: ' +
' The PassTicket generation was a success but ' +
' unfortunately there was also an error while trying to ' +
' print it using the TSO/E PUTLINE macro. ' +

PTKN019 . ALARM=YES . TYPE=ACTION . WINDOW=LR
' Unknown error. Please contact your systems programmer. ' +
' System programmers action: correct the error and rerun the job. ' +

PTKN02 Member

PTKN020 . ALARM=YES . TYPE=WARNING . WINDOW=LR
' The PassTicket generated for user &USER and VTAM application &APPL ' +
' is &PKT. . This will be valid for 10 minutes and can only be used ' +
' once. ' +

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E-mail alerts

Our e-mail alert service will notify you when new issues of *RACF Update* have been placed on our Web site. If you'd like to sign up, go to <http://www.xephon.com/racf> and click the 'Receive an e-mail alert' link.

How RACF handles passwords

This article discusses the issues surrounding password authentication. It starts by reviewing the principles behind authentication and encryption, and then develops this theme in order to arrive at a clearer understanding of how RACF handles passwords. Finally, it compares RACF and Windows NT, to bring out the main cross-platform issues associated with synchronizing RACF and NT passwords in a single-sign-on scenario.

Note that whenever I use the terms 'RACF' and 'Windows' in this article, they can be taken to mean RACF or the more recent OS/390 Security Server, and Windows NT/2000/XP.

PASSWORD AUTHENTICATION

The topic of passwords is clouded in mystique and misconceptions, as can be seen by the questions and discussions that periodically arise on racf-l¹, the Internet discussion list for RACF. So first, let's see if we can sort out some basics.

Authentication is a means of being able to arrive at some level of confidence that the person you are communicating with is indeed who they say they are. Humans use many ways to authenticate people – for example, the sound of the person's voice, visual recognition, etc. Computers are much more limited in the ways in which they can try to authenticate who they are communicating with. Typically, they use a combination of userid and password, and, increasingly these days, some form of software or hardware token. Tokens are generally used as a means of implementing one-time passwords (see below). There are three primary opportunities to attack passwords:

- When they are entered
- When they are transmitted
- When they are stored.

Entering passwords

When a password is entered, you're basically at the mercy of the device through which you're entering it. If it's a secure entry device, it should be secure. If it's a general-purpose PC which could potentially have a key-logger installed, it's not so secure.

Transmitting passwords

Two strategies can be used to protect against the interception of transmitted passwords. One is to encrypt the communications channel (eg SSL, tn3270e), and the other is to use a one-time password so that, even if someone does intercept the password, it doesn't matter because they can't re-use it.

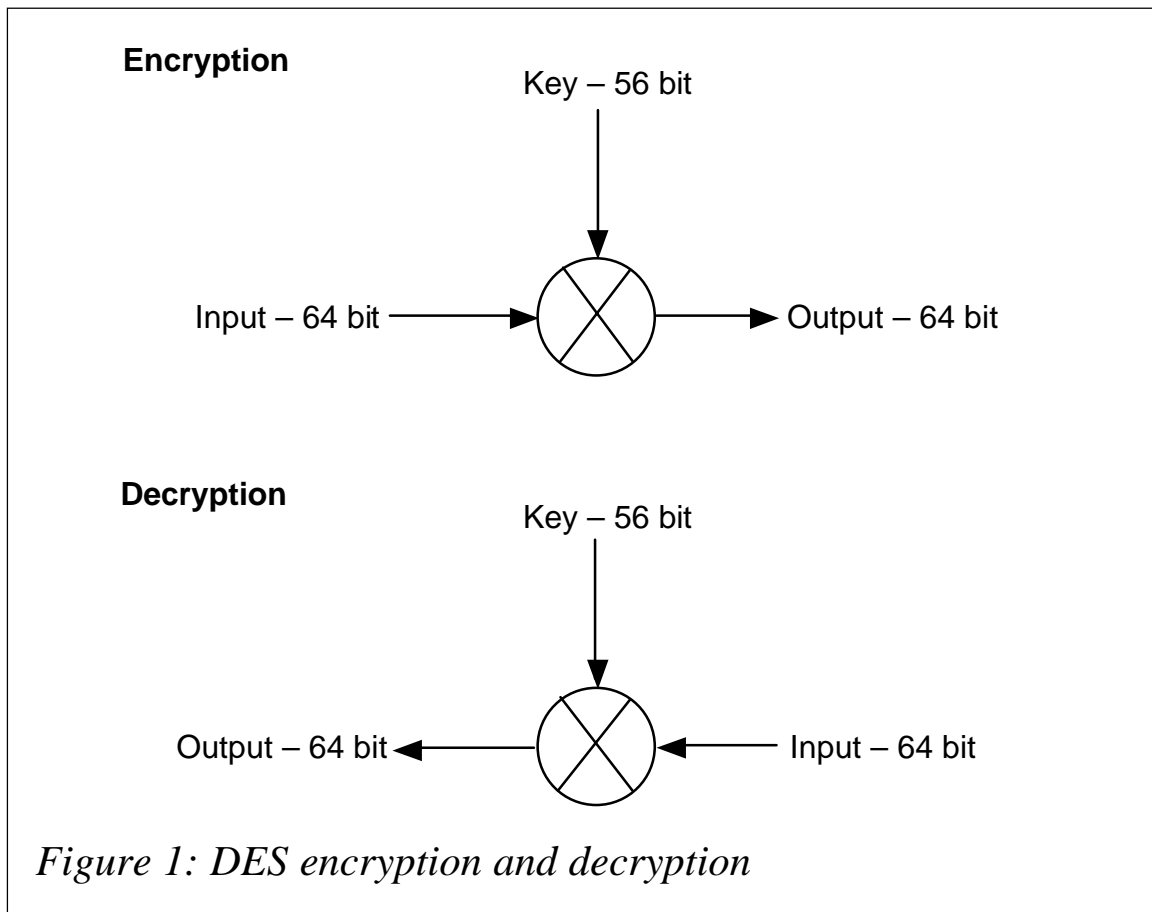
Storing passwords

Last, and I think generally the most important, is the question of how passwords are stored. This is particularly important because an attack on a password when it's entered or transmitted is an attack on just one password. Attacking the store, by contrast, means being able to attack the entire user population. This makes it a very attractive target.

Passwords should never be stored on a system in such a way that they can be retrieved. Instead, they should be encrypted using a one-way function to arrive at an encrypted value. The sign-on process should then go through exactly the same process and compare the encrypted result with the stored encrypted value to see if they match.

This is fundamental, and most operating systems do a reasonable job, but many applications do not!

Don't forget that although I'm concentrating on the stored repository of encrypted passwords, equally vulnerable is any mechanism used to store them, eg the ICHPWX01 password exit. If someone can install their own code in here, it's pretty much 'game over' anyway.



THEORY OF ENCRYPTION

The Data Encryption Standard² (DES) has been around since 1977³ and has been the base building block for many systems, including RACF. It's relatively simple in theory, even if the precise implementation can get tricky.

From the outside, it works like a black box that has two inputs and one output. The inputs are a 64-bit block of data and a 56-bit key. The black box uses the key to encrypt the data into a 64-bit output data block. To decrypt, you do the inverse – that is, you use the same key but reverse the algorithm (see Figure 1).

DES is normally described starting from the outside and then working in. Here, however, I'll focus simply on the core mechanism used by DES, and then move outwards enough to set it in context. This should, hopefully, be enough to give an appreciation of its properties.

At the core of the DES algorithm are a series of eight 'S' boxes

S1															
14	4	13	1	2	15	11	8	3	10	6	12	5	9	0	7
0	15	7	4	14	2	13	1	10	6	12	11	9	5	3	8
4	11	4	8	13	6	2	11	15	12	9	7	3	10	5	0
15	12	8	2	4	9	1	7	5	11	3	14	10	0	6	13

Figure 2: S box in detail

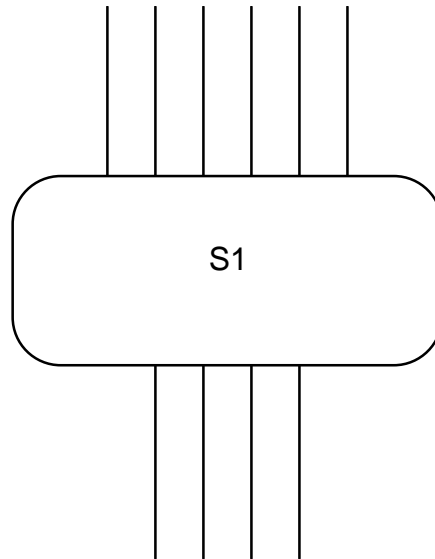
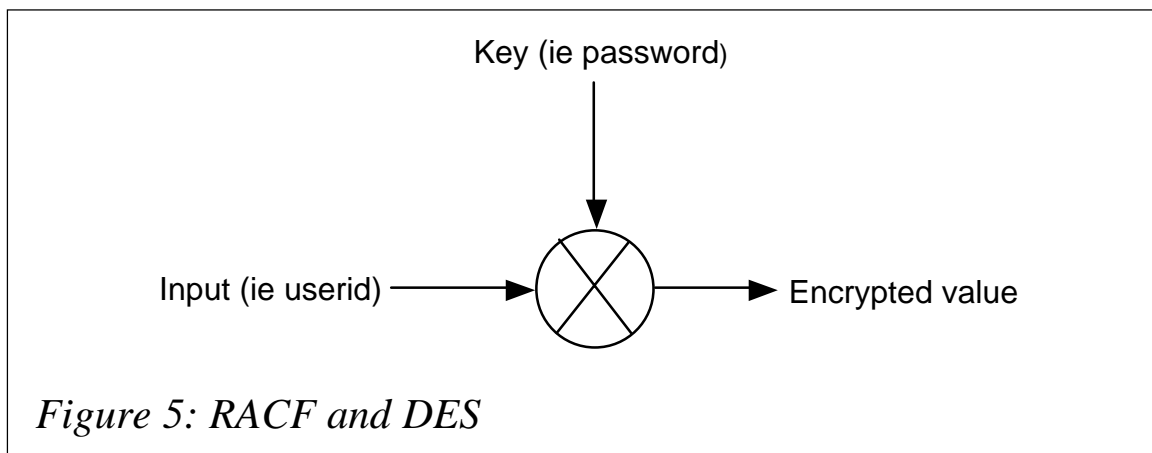
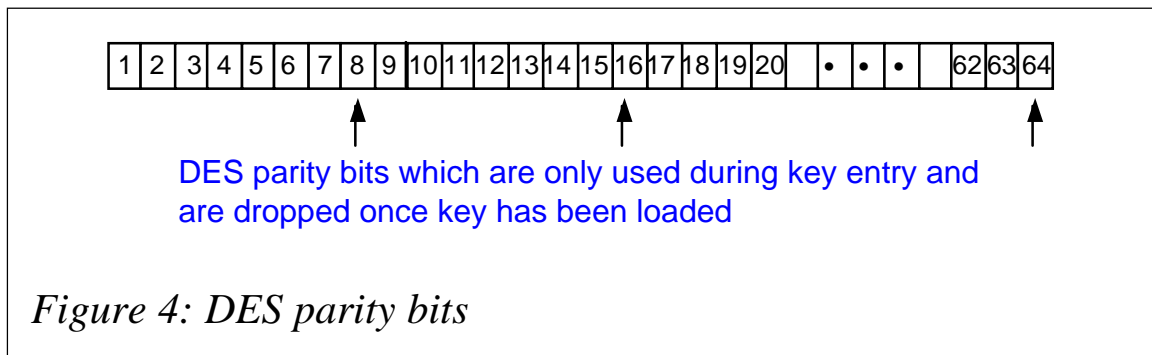


Figure 3: S box concept

(see Figure 2). The 'S' box operates by acting as a look-up table, where the input value is a 6-bit number. Two of the 6 bits are used to decide the row, and the remaining 4 bits are used to determine the column. This then gives an output number between 0 and 15, ie a 4-bit number. This is generally represented as shown in Figure 3.

The other seven 'S' boxes are similar tables, but with the values rearranged to give a different set of look-up values, otherwise just a variation on a theme.

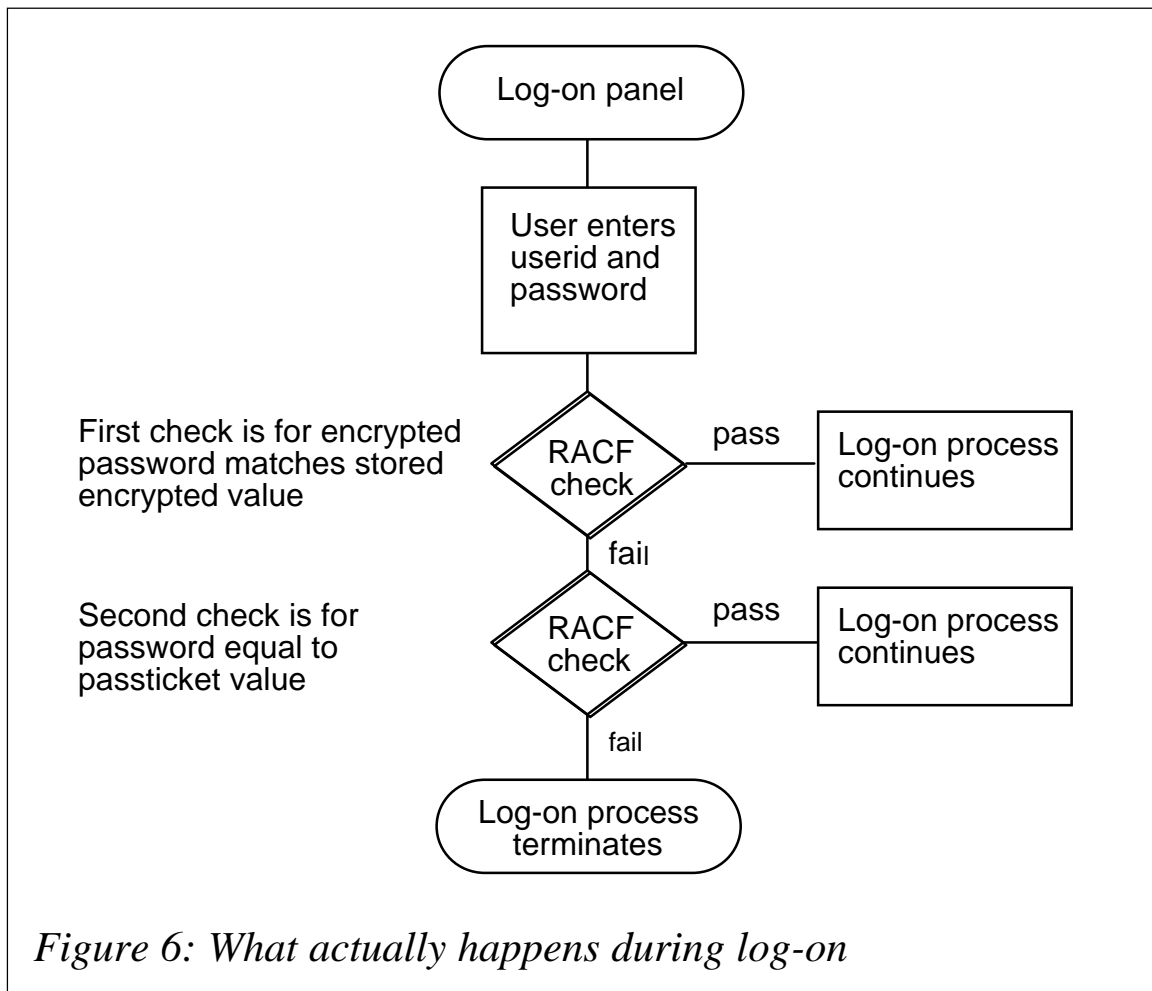
If you understand the 'S' box concept, you almost understand DES, because DES simply takes this principle and places it in a series of nested loops, so that to try to work it back becomes computationally infeasible – that is, it will require as many



computations as trying every permutation of the 56-bit key. So even if you know the userid and the encrypted value, the only way to attack it is to use either a dictionary attack or brute force.

In cryptographic terms, the whole point of DES is to protect the key. A cryptographic attack would typically involve obtaining some input data along with some associated encrypted output data, and then trying to deduce the key. DES can either be used as a method of encrypting data so that it can be decrypted at a later time, or as a one-way function – ie throwing away the key once it has been used to encrypt some input data.

Incidentally, in case you're wondering why DES uses a 56-bit key rather than a 64-bit key, allow me to explain. It all dates from the time when keys were manually input to secure devices through keyboards or keypads. They would typically take a 64-bit value as input, but the device would then break it down into 8-bit blocks where each block contains seven data bits and one parity bit (see Figure 4). As this was generally keyed as 16 hexadecimal digits, this gave a rudimentary level of error checking to detect keying errors.



RACF PASSWORDS

I'm pleased to say that RACF uses DES as a one-way function, as shown in Figure 5. Note that although IBM uses DES as per the standard, it has to take an 8-character password and convert it into a 56-bit key. Understanding how it does this requires us to look down a level at the binary bits involved.

As we've seen, in DES the right-most bit of each byte is lost as a parity bit. Now consider the EBCDIC characters, hex values, and bit patterns involved:

'A' = C1 = 11000001
 'B' = C2 = 11000010
 ...
 '8' = F8 = 11111000
 '9' = F9 = 11111001

plus three special country characters. In other words, the left-

most bit is almost always a 1, and hence virtually redundant, leaving 7 useful bits. Yet again, there's no rocket science here, just good sound logic.

RACF passwords are stored in the RACF database datasets as identified by the RVAR Y LIST command, so it's not difficult for any user to find out where the encrypted passwords are. These datasets must therefore be protected against unauthorized reading.

Figure 6 shows what actually happens during log-on. PassTicket is techno jargon for an IBM proprietary method for one-time passwords using a software implementation of a token – that is, software used to generate a one-time password based on a secret key configured into the software, the userid, and, of course, a function of the date and time. If you want to delve any deeper, I'd recommend Thierry Falissard's *The RACF PassTicket Page*⁴ as an excellent starting point.

PassTickets are of particular benefit if connecting to a system across an untrusted network where there's a possibility of someone trying to 'sniff' userids and passwords. However, this advantage has to be weighed against the fact that the client system needs to have user credentials stored locally. The risk here becomes a combination of the accessibility of the client system and the manner in which the credentials are stored – eg in the clear or encrypted.

By now it should be becoming clear that there's no rocket science involved here: anyone who can go and look up the DES on the Internet and apply it to RACF can carry out a dictionary attack.

Several RACF password crackers are now freely available, in addition to the genuine security administration and audit tools which contain password crackers. The earliest cracker was from Kurt Meiser (now marketed by Peter Goldis⁵). Then came my CRACF⁶, followed by another from Thierry Falissard, both available over the Internet. Most recent is an evolution from CRACF called WEAKWORD⁷. CRACF displayed any cracked

passwords but was very restrictive in what it tried to crack. WEAKWORD, on the other hand, doesn't display the password but just flags it as weak; however, it does allow a dictionary to be defined. Interestingly, I received more requests for a version which didn't display the cracked passwords than I did for a full-blown cracking version.

There are several ways in which we can mitigate the risk of attack to RACF passwords:

- Apply password rules or use the password exit, to reduce the likelihood of guessable passwords.
- Ensure password history is used to prevent simple recycling of passwords.
- Ensure passwords are not static and are changed periodically.
- Protect the encrypted password storage from unauthorized read access.

PASSWORD SYNCHRONIZATION

As we've seen, RACF doesn't actually store encrypted passwords, but rather encrypted userids, where the password was used as the key and then discarded. This means that if two userids both have the same password, the encrypted values stored in the RACF database will be different. However, this is not true of all systems and is certainly not true of Windows NT.

Windows also uses a much less computationally intensive algorithm for encrypting passwords, making it a very attractive target for hackers. Put simply, it's quicker to carry out an attack on an encrypted Windows password than on an encrypted RACF password, and, more importantly, it takes virtually the same time to attack one Windows password as it does to attack an entire Windows population of users. Compare this with the much older RACF, where every password must be individually attacked.

I don't mean to delve too deep into Windows passwords, but I do feel it's worth pointing out how they're processed/stored. In particular, I'm referring to the fact that they're split into blocks of seven characters before they're encrypted and stored. This means that a 14-character password is not much better than a seven-character one. If you can brute-force-attack a seven-character password, you can brute-force any length of Windows password.

What's more, if your password is, say, nine characters long, then the attack analyses the second (two-character) block first, and will crack that with very little effort. This then leaves the initial seven characters, along with the additional clue of knowing what the last two characters are. Worse still is if your password is, say, a seven-character word, followed by two numbers. My advice with Windows passwords is to use seven-character passwords and be sure to include at least one punctuation character⁸.

Incidentally, Windows also has the equivalent of the ICHPWX01 password exit. Any dynamic link library located in the %systemroot%\system32 directory and referenced in the HKEY_LOCAL_MACHINE\SYSTEM\CurrentControlSet\Control\Lsa\Notification Packages registry key will be called during password change. It's very much the equivalent of the RACF exit, only the RACF exit is easier to control as it's only in one place. The Windows exit relates only to the local password database – ie if the exit is on a workstation it relates only to the local passwords held on that workstation. Likewise, if the exit is on a Primary Domain Controller, it relates to every password stored locally on that domain controller.

Microsoft supplies a sample exit called PASSFILT.DLL9 for Windows NT which applies its suggested rules for strong password validation – ie there are no configurable options, and you either use it or write your own. The functionality of PASSFILT.DLL has been incorporated into Windows 2000 and XP.

I can't talk about Windows passwords without the obligatory

reference to the *de facto* industry standard cracker, L0phtcrack¹⁰. This is a very powerful and impressive tool which all information security professionals should be aware of. If you've never seen it, go and try it – you'll learn something!

Most recently, in July 2003, there has been some very interesting work¹¹ coming out of LASEC¹² (Security and Cryptography Laboratory) which exists within EPFL¹³ (École Polytechnique Fédérale de Lausanne). It has developed a very fast password cracker called Advanced Instant NT Password Cracker. To demonstrate how fast, it provided an on-line interface and invited visitors to submit encrypted passwords for cracking. They managed to crack 1,845 passwords with an average crack time of 7.7 seconds!

I hope I've managed to convince those who weren't already convinced that there are some very significant differences in the strengths of design between the ways different operating systems handle passwords. This makes me very wary about the potential of synchronizing systems such as a password reset on a Windows platform, which can be propagated to an associated RACF userid, or vice versa.

CONCLUSIONS

In summing up, I'd like to start by quoting from Bruce Schneier, an academic cryptographer cum real-world security practitioner¹⁴: “You can't memorize good enough passwords any more, so don't bother. Create long, random passwords, and write them down. Store them in your wallet, or in a program like Password Safe. Guard them as you would your cash. Don't let Web browsers store passwords for you. Don't transmit passwords (or PINs) in unencrypted e-mail and Web forms. Assume that all PINs can be easily broken, and plan accordingly.”

Password Safe¹⁵ is a freeware utility which started life at Counterpane under the direction of Bruce Schneier. It's been around for many years unchanged. However, it's now being actively developed as a SourceForge Open Source Project¹⁶. It's a small application dedicated to storing userids and passwords securely.

Let me say that I agree whole-heartedly with Bruce Schneier's statement. I for one use Password Safe and actively encourage its use wherever and whenever I can.

The only caveat with Password Safe is that it is only as secure as the password used to secure it in the first place, and there is a password safe cracker¹⁷ by Joe Smith freely available on the Internet which can be used to carry out a dictionary attack on Password Safe. This means that the initial password used must not be something which is ever likely to appear in any password hacking dictionary.

Because Password Safe is a small, simply run program without requiring libraries etc, it lends itself to being kept on USB removable media, giving an added level of physical security.

Finally, just to recap on the points touched upon earlier, here is my checklist:

- Use Password Safe or an equivalent.
- Use password expiry to enforce regular password changing.
- Use setropts rules to enforce improved password quality.
- Use password history in conjunction with setropts password rules.
- Ensure that the RACF database datasets, primary and back-up, have fully qualified dataset profiles that have no unnecessary read access.
- If authenticating across an insecure network, consider using PassTickets as an alternative to conventional passwords.
- Keep passwords different across different platforms, at least for privileged users.
- Know about any password exits in use within the environment, whatever the platform.
- Use seven-character Windows passwords with at least one punctuation character.

REFERENCES

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- 3 Applied Cryptography Second Edition: protocols, algorithms, and source code in C / Bruce Schneier
- 4 <http://os390-mvs.hypermart.net/passtkt.htm>
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RACF in focus – understanding OS/390 Unix security

‘RACF in focus’ is a regular column focusing on a specific RACF topic. Here, we examine RACF security for OS/390 Unix.

OS/390 Unix was initially known as OpenEdition MVS, but that name is no longer used. It is now commonly referred to as OS/390 Unix System Services (USS), or simply OS/390 Unix.

Security checking for OS/390 Unix is done in RACF, and in order to use Unix services a user must have a RACF userid. The RACF security administrator therefore needs to understand RACF security as it relates to OS/390 Unix, and implement at least portions of it to protect Unix resources.

HOW OS/390 UNIX SECURITY WORKS

Unix data is stored in files whose structure is more like the one used in Windows-based PCs than the datasets used in the MVS environment. This file system is known as the Hierarchical File System (HFS). Security information, consisting of flags, is kept within the file system in File Security Packets (FSP). Within the FSP, the UID represents the file owner, and the GID represents the group owning the file. Access is given based on UID and GID information stored in the FSP.

On the RACF side, there’s an OMVS segment in the RACF userid profile (much like the CICS and TSO segments) that defines various Unix attributes for the user, including the UID, and an OMVS segment in the group profile that defines the GID. Before a RACF user can use Unix services, he or she must have an OMVS segment (or inherit a default, as described below). The user also needs to be connected to a group with an OMVS segment (or inherit a default).

When Unix resources are accessed, security checking is done by calling RACF. RACF compares the UID and GID values of the Unix resource to the UID and GID values assigned to the

user. The UID and GID values assigned or inherited by a user therefore play an important part in the Unix world.

PLANNING FOR OS/390 UNIX SECURITY

If your installation is going to deploy Unix applications in the IBM mainframe environment, you'll obviously need to plan for and implement RACF security for OS/390 Unix.

However, there are a number of other reasons why you might need to implement some portions of RACF/Unix security, even if you're not actually implementing Unix applications. For example, if you want to do file transfers between the mainframe and other platforms using ftp, you'll need OS/390 Unix services, and, therefore, RACF security. TCP/IP and LDAP implementations are other examples that require OS/390 Unix and therefore RACF security.

It's important to have a plan that defines basic RACF security for a Unix environment, so you'll be well positioned to exploit its features more fully if and when required. Amongst other things, you'll need to define policies and procedures to administer the Unix environment, and decide who is to be allowed to use OS/390 Unix, what UIDs and GIDs to assign, how you'll track assigned values, how you'll monitor and audit the UnixC environment, and so on. Without a plan, you may end up having undesirable results. For example, ftp usage is quite common these days, and RACF administrators are often asked to assign OMVS segments to users wanting to use ftp processes. Without a plan, you may assign OMVS segments without paying much attention to the UID and GID values. This may have negative implications in the future, and some of the work will need to be re-done.

UNIQUE UIDS AND GIDS RECOMMENDED

Although UIDs and GIDs can be shared among users and groups, this isn't recommended. Sharing doesn't really make sense because file protection in the HFS is done at the UID and GID level, and you lose accountability.

Ideally, each user should have a unique UID assigned, and each group a unique GID. The only exception is UID(0), which gives the user 'superuser' powers. You may have more than one user who needs this.

The RACF class UnixMAP keeps tabs on who is assigned which UIDs and GIDs. It can be used to ensure that RACF users and groups have unique UIDs and GIDs. It's best to activate this class before you start assigning UIDs and GIDs at your installation.

If you can't do this, you'll have to take the action described in the *OS/390 Security Server (RACF) System Administrator's Guide*. This involves running a REXX EXEC to 'populate' the UnixMAP class from the RACF database and then activating the UnixMAP class to keep tabs on future changes to UIDs and GIDs.

If the UnixMAP class is active, you can query it to find out who is using a particular UID or GID, as follows:

```
RL Uni xMAP G555 al l
```

It's important to note that profiles in the UnixMAP class are auto created – you don't need to maintain the access list.

You may not want to assign OMVS segments to all userids and all groups, even though they require Unix services. In this case, you can specify defaults to be used by individuals who don't have any OMVS segments. This is done via the FACILITY class profile BPX.DEFAULT.USER. It contains, in the appldata field, the default RACF userid and default RACF groupid, to be used for OMVS segment look-up purposes. You have to define the default userid and the default groupid to RACF. Note that the defaults are used only in cases where no OMVS segment is found for the user. If an OMVS segment is found, information from that segment is used, not from the default userid's segment.

SUPERUSER POWERS

In the Unix world, UID(0) is used to denote a superuser. This is a very powerful attribute, and can be compared to someone having the RACF attributes OPERATIONS, SPECIAL, and

AUDITOR all at once. Note that the RACF attributes OPERATIONS, SPECIAL, and AUDITOR give no special powers in the Unix world; similarly UID(0) gives no special powers for MVS purposes.

If you have the UnixMAP class active, you can see how many users have UID(0) assigned by entering the following command:

```
RL Uni xMAP UO ALL
```

Since superuser is such a powerful authority, you may not wish to give this to many people. Some users may need only a subset of the superuser powers. There is a RACF class called UnixPRIV that you can use to specify more granular levels of special powers.

AUDITING OS/390 UNIX

The following classes are available for auditing OS/390 Unix: DIRACC, DIRSRCH, FSOBJ, FSSEC, IPCOBJ, PROCACT, and PROCESS. These classes don't need to be activated, nor do you need to create profiles within these classes.

In order to audit OS/390 Unix, you need to specify your audit options using either of the following commands for the above classes:

- SETROPTS LOGOPTIONS(CLASSNAME)
- SETROPTS AUDIT(CLASSNAME)

SUMMARY

This article has looked only very briefly at OS/390 Unix security, and contains enough information to get you started. If your installation is developing OS/390 Unix applications, you'll need to do further research into the workings of OS/390 Unix. Finer control mechanisms are available in OS/390 to address more complex Unix configurations.

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RACF news

Critical Path has announced its Critical Path Password Management application for centrally administering passwords across systems and applications. The integrated software provides self-service resets of forgotten passwords, centralized definition and enforcement of password policies, dynamic password synchronization across systems for reduced sign-on, and auditing of all password change activities. It enables passwords for various user applications, databases, and operating systems to be administered centrally. Users can set their own passwords via a Web-based interface or through existing systems such as RACF or Windows.

URL: <http://www.criticalpath.net/solutions/enterprise/passwordManagement/>

* * *

OpenNetwork Technologies has announced the Universal Identity Platform (Universal IdP), designed to take advantage of Microsoft technology and to extend its value as an identity infrastructure to J2EE and mainframe environments. The software centralizes and unifies the management of identities and security policies, secures access to protected resources, and delivers automated workflow and provisioning. Specifically, it integrates Microsoft Identity Integration Server (MIIS) with mainframe systems such as RACF, ACF2, and TopSecret.

URL: http://www.opennetwork.com/news/press/2003/2003-07-02_UIIdP.php

* * *

IBM has announced Tivoli Workload Scheduler Version 8.2, designed to help reduce the complexity of managing the workload on mainframes and open systems and automating many operator activities. Enhancements include improved security through the addition of SSL-based authentication and encryption.

URL: <http://www-3.ibm.com/software/tivoli/products/scheduler-apps/>

* * *

e-Security has announced the release of e-Security Version 4, with new functionality for managing enterprise security, including enhanced usability, incident management, performance, and correlation capabilities.

URL: http://www.esecurityinc.com/Company/Press_Releases/Dynamic.asp?PR_ID=26

* * *

Computer Associates and SteelCloud have announced an agreement under which the companies will deliver a family of hardened, ready-to-deploy enterprise-class security appliances based on CA's eTrust family of security solutions.

URLs:
<http://www3.ca.com/press/PressRelease.asp?CID=45782>
<http://www.steelcloud.com/appliances/default.asp>

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